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IN TRODUCTION !

Digital signal processing to an area of science and engineering that has developed rapidly over the past he years. This rapid development & a result of the significant advances in digital computer technology and integrated - circuit febrication. The rapid developments in IC technology starting with Medium Scale Integration (HSI) and progressing to Large Scale Integration (HSI) and progressing large Scale Integration (HSI) and progressing development of powerful, smaller, feaster and cheaper digital computers and special-perpose digital hardware.

T -[1-6]

These digital circuits have made it possible to construct highly sophiliticated digital systems capable of performing complex digital signal processing functions. May of the signal processing tasks that where conventionally many of the signal processing tasks that where conventionally performed by analog means are realized togy by less expensive and often more reliable digital heroderane.

Processor

SIGNALS !-

A signal & defined as any physical quantity that varies with time, space or any other independent variable tors variables. It & described mathematically as a function of one or more independent variables. Ex: Si(t) = 5t Solt) = 20t²

 $S(x,y) = 3x + 2xy + 10y^{2}$

EnggTree.com @ System'-A system to defined as a physical (or software realizations) device, that performs an operation on a signal Ex: A filter used to reduce the noise and interference correpting a desired information bearing Bignal & called a system. SIGNAL PROCESSING!-In general, a system is characterised by the type of operation that it performe on the signal. Such operations are usually referred to as signal processing. the Lorge Scale Indigrad BASIC ELEMENTS OF DIGITAL SIGNAL PROCESSING Most of the signals encountered in science and engineering are analog in nature. To process these analog signals, appropriate analog systems are used. The illustration of analog processing system & shown below. party made by analog made Analog inpout Signal Drocessor Signal signal as any physical quants Figure 1.1 Analog signal processing system described mathematical that yarde with i lines Processing of analog signals may include changing their characterentics cors extracting their desired information. Here both the input and output are in analog form

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Digital signal processing provides an alternative method for processing the analog 3 signal as shown below in figure 1.2. talton. DIA Digital Analog Analog A/D Digital converter Output rignal signal Signal Processor signal Figure 1.2 Digital Signal Processing System To process analog signal in digital processor, an interferce called ALD convertor & to be used. The output of the A/D converter & a digital signal that is appropriate as an input to the digital processor. The digital signal processor may be a large programmable digital computer or a small microprocusor programmed to perform the desired operations on the input signal. In applications like speech communications where the digital output from the processor & its be given to the user in analog form, another introquee called D/A converter is used. ADVANTAGES OF DIGITAL OVER ANALOG SIGNAL PROCESSING A digital programmable system allows filexibility in reconfiguring the digital signal processing (i) operations simply by changing the program. (ii) Digital systems are more accurate (iii) Digital signals can easily be stored on magnetic modia without loss of signal ficklity.

EnggTree.com (iv) Inplementation of more sophisticated signal processing algorithms as easy in digital signal processing. (V) Digital implementation of signal processing systems & cheaper than its onalog counterpart, Bralo Applications !-(i) Speech processing & signal transmossion on telephone channels. (ii) I mage processing and france lision (iii) Selemology and geophysics (iv) In oil exploration (V) The detection of nuclear explosions (vis processing of signals received from outer programmabale digital compater or a stall micropoor (vii) Vast variety of other applications. Limitation: input signal. (i) Speed of operation of ALD conventures and oligital signal processors. (i) Signals having wide range of bandwidth require fast sampling rate Als conversions and fast digital signal processors. A digital programmable agatem allows of leaders (1) is recentizerize the digital algori proceeding operations simply by changing the program

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(3)

DISCRETE FOURIER TRANSFORM

Descrete Signals and Systems - A Reetew Introduction to DFT. - properties of DFT. Circular convolution - Fidtering methods based on DFT - FFT algorithms - Decimation in time Algorithms, Decimation in frequency. Algorithms - Use of FFT in dinear filtering.

DISCRETE SIGNALS AXID SYSTEMS _ A REVIEW T - (41-67) Signals that are defined only at certain specific values of time are called discrete-time signals. These time instants ned not be equidistant, but in practice they are usually taken at equally spaced intervals for computational convenience and mathematical tractebility. A déscrete-time signal having a set of déscrete values & called a digital signal. ie, A signal which & ducrete in strue and amplitude & called digital signal. Elementary Descrete - Time Signals :-O Unit sample sequence los unit impalse $\sigma(n) = 2 \circ n \neq 0$ roomer - come

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Figure 1.3 Unit impulse signal

1012 ... n

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(1) Unit step signal

$$u(n) = \int_{0}^{1} (n + 20)$$

$$(u(n) = \int_{0}^{1} (n + 2)$$

$$Fgure 1 + unit step signal
(2) Unit Eamp signal
$$u_{y}(n) = \int_{0}^{n} n + 20$$

$$(u_{y}(n) = \int_{0}^{n} n + 20)$$

$$(u_{y}(n) = \int_{0}^{n} n + 20$$

$$(u_{y}(n) = \int_{0}^{n} n + 20)$$

$$Fgure 1 + 5 Unit ramp signal
(2) Exponential signal
$$x(n) = \alpha^{n} \quad \text{for all } n$$

$$if a = u \quad read \quad \text{fiers } valued \quad \text{fiers } u$$

$$u_{y}(n) = x^{n} e^{2n}$$

$$u_{y}(n) = x^{n} e^{2n}$$$$$$

EnggTree.com 0 x(n) rem Caci a>1 · (a) = (a) x - 1-20 x(m) & called ontrapponetric (A) x -- tota :2H 3T2H2 33MLCI Classification of Discrete - Time Signals:-* Energy signals and Power signals Energy signals $E = \frac{5}{2} |r(n)|^2$ If E & finite, P=0 $\sum_{i=1}^{N} |z(n)|^2$ Power Signale p= lim _____ N+1 D=-N = Lim 1 EN N-30 2N+1 EN where EN + Signal energy over -NENEN If E & o, the ave power may be finite (00) infinite. * Periodic Signals and appriodic signals A signal ress & periodic with period N(N>0) if and only if $x(n+x) \equiv x(n) + n$ If there is no value of N that satisfies the above equation, the signal is called nonperiodic (or) aperiodic.

EnggTree.com Symmetric (over) and antisymmetric codd) signals! A real valued signal xind . to signitude if $\chi(-n) = \chi(n)$. A signal x(s) is called antisymmetric if 2(-n) = -2(n)· · · · · · · · DISCRETE TINTE SYSTEMS! A déscrete-tonie systèrs & a device cos algorithm that operates on a descrete time signal

A discrete-time system & a device (o) algorithm that operates on a discrete-time signal called the input (or) excitations, according to some well defined rule, to produce another another discrete time signal called the output (or) response of the system.

Classification of Dicrete-Time Systems:

* Static Versus dynamic Systems:-

A déscrete-tenu system & called statio con memoryless if its output at any instant in depends at most on the input sample at the same time, but not on past (or) future sample of the input.

A system in said to be dynamic, if it

has a memory.

EnggTree.com a & Time - invariant versus time - variant systems A released system T is time invariant Low shift invariant if and only if $x(n) \xrightarrow{T} y(n) \xrightarrow{I-1} x$ implies that r(n-k) -> y(n-k) for every input xing and every line shift k. A Linear Versus XIon-linear: A system & linear if and only if $T\left[a, x, (n) + a_2 x_2 (n)\right] = a, T\left[x, (n)\right] + a_2 T\left[x, (n)\right]$ for any arbitrary input sequence x, (1) and x.0) and any arbitrary constant a, and as. * Causal Versous Non causal systems A system & said to be causal if the output of the system at any time & depende only on present and past inputs but not depende on fature inputs. i y(n) = F[x(n), x(n-1), x(n-2)...]* Stable Versus Onstable système An arbitrary relaxed system & said to be bounded input - bounded output (BIBO) stable if and only if every bounded input produces a bounded patput. $z x(0) e^{-\frac{1}{2}} + x(0) e^{-\frac{1}{2}} + x(0) e^{-\frac{1}{2}} + x(0) e^{-\frac{1}{2}}$

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Discret to Fourier Transform
$$Ti - (hTh - hTh)$$

The Blocalt Fourier Transform (DFT) of
a Agred relia & granby
 $x(k) = \sum_{n=0}^{N-1} x(n) e^{-j\frac{n}{N}} keo, 1, 2... M_{rh}$
 $x(k) = \sum_{n=0}^{N-1} x(n) e^{-j\frac{n}{N}} keo, 1, 2... M_{rh}$
The inverse DFT (EDFT) & granby
 $x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) \perp \frac{j\frac{n}{N}}{N}, neo, 1, ... M_{rh}$
DFT & the frequency domain sampling of
a signed $x(n)$ of length N .
O Compute the DFT of the sequence $x(n) = \int 1... M_{rh}$
 $(AHI-SOIT-RIS)$
Cliven $x(n) = \int 1... -1, 1... -1 \int (AHI-SOIT-RIS)$
 $Chiren $x(n) = \int 1... -1, 1... -1 \int (AHI-SOIT-RIS)$
 $X(k) = \sum_{n=0}^{N-1} x(n) e^{-j\frac{n}{N}} koo, 1, ... M_{rh}$
 $x(k) = \sum_{n=0}^{N-1} x(n) e^{-j\frac{n}{N}} koo, 1, ... M_{rh}$
 $x(o) = \sum_{n=0}^{3} x(o) e^{-j\frac{n}{N}} koo, 1, ... M_{rh}$
 $x(o) = \sum_{n=0}^{3} x(o) e^{-j\frac{n}{N}} koo, 1, ... M_{rh}$
 $x(o) = \sum_{n=0}^{3} x(o) e^{-j\frac{n}{N}} koo, 1, ... M_{rh}$
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 $x(o) = \sum_{n=0}^{3} x(o) e^{-j\frac{n}{N}} koo, 1, ... M_{rh}$
 $x(o) = \sum_{n=0}^{3} x(o) e^{-j\frac{n}{N}} koo, 1, ... M_{rh}$$

$$\begin{array}{l} & & & \text{EnggTree.com} \\ & & & x(5) \neq (0(45) + (-5)(45) + (5)(6)(4) + (5)(6)(4) + (-3)(4) \\ & & X^{(1)} = 1 + (-1)(6-1) + 1(-1) - 1((6+1)) \\ & = 1 + (-1)(6-1) + 1((-1)) - 1((6+1)) \\ & & X^{(1)} = 3 \\ & & X^{(2)} = 3 \\ & & X^{(2)} = 3 \\ & & X^{(3)} = -\frac{1}{2} \frac{\pi^{(3)}}{\pi^{(3)}} \\ & & = 1 - 1(-5) + 1(15) - 1(-1) \\ & & = 1 + (1 + 1 + 1) \\ & & = 1 + (1 + 1 + 1) \\ & & X^{(1)} = 3 \\ & & X^{(2)} = 3 \\ & & X^{(3)} = 3 \\ &$$

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Relationship of the DFT to other transforms
(D) Relationship to the Fourier series coefficients
of a periodic sequence
$$[z_p(n)]$$
 with
fundamental period kl can be represented in
a Fourier series of the form
 $z_p(n) = \sum_{k=0}^{N-1} c_k e^{j\frac{2\pi}{N}} - det nco$
where Fourier series co-afficients are given by
 $c_k = \frac{1}{N} \sum_{n=0}^{N-1} z_p(n) e^{-j\frac{2\pi}{N}} kz_{0,l,nks}$
Comparing these equations with DFT and
 $TDFT = equations,$
 $x(k) = MCk$ and
 $z_p(n) = \frac{1}{N} \sum_{k=0}^{N-1} x(k) e^{-j\frac{2\pi}{N}} hz_{0,l,nks}$
 $k=0,l_{0,nks}$
 $k=0,l_{0,nks}$
 $k=0,l_{0,nks}$
 $k=0,l_{0,nks}$
 $k=0,l_{0,nks}$
 $k=0,l_{0,nks}$
 $h=1$ $x(k) e^{-j\frac{2\pi}{N}} hz_{0,l_{0,nks}}$
 $h=0,l_{0,nks}$
 $h=1$ $x(k) e^{-j\frac{2\pi}{N}} hz_{0,l_{0,nks}}$
 $h=0,l_{0,nks}$
 $h=0,l_{0,nks}$

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Pelationship to Fourier Transform of an
apertodic sequence (DTFT)
If Z(N) & an approduce finite energy
requere with Fourier Transform X(W) which is
requere with Fourier Transform X(W) which is
requere with Fourier Transform X(W) which is
requere to
$$M = equally spaced firstundies
Wh = $2\pi k$, $k = 0, 1, \dots, N-1$ the spectral
 $Wh = 2\pi k$, $k = 0, 1, \dots, N-1$ the spectral
 $Wh = 2\pi k$, $k = 0, 1, \dots, N-1$ the spectral
 $Wh = 2\pi k$, $k = 0, 1, \dots, N-1$ the spectral
 $W(W) = \chi(W) / W = 2\pi k$
 $M = \frac{5}{N} - \chi(N) = \frac{1}{2\pi n k}$ ($k \ge 0, 1, \dots, N-1$
 $h = \infty$
are the DFT coefficients of the periodic
 $xequere of period N, given by
 $w_p(n) = \frac{5}{N} - \chi(n - LN)$
 $U = -\infty$
 $M = 4 transform of a sequene $\chi(N)$ G
 $\chi(W) = \frac{5}{n_{m-\infty}} - \chi(n - LN)$
 $Wh k = 2\pi k hat include the work circle.
If $\chi(z) = \frac{5}{n_{m-\infty}} - 2^{(n)} z^{-n}$
 $with k = 4 transform of a sequene $\chi(N)$ G
 $\chi(k) = \frac{5}{n_{m-\infty}} - 2^{(n)} z^{-n}$
 $with k = 1 the work circle $z_k = e^{\frac{1}{2\pi k}} dkm$
 $\chi(k) = \chi(z) / z = a \frac{1}{2\pi k}$ $k = 0, 1, \dots, N - r$
 $\chi(k) = \frac{5}{n_{m-\infty}} - x(n) e^{-\frac{1}{2\pi k}}$ $k = 0, 1, \dots, N - r$
 $\chi(k) = \frac{5}{n_{m-\infty}} - x(n) e^{-\frac{1}{2\pi k}}$ $k = 0, 1, \dots, N - r$
 $\chi(k) = \frac{5}{n_{m-\infty}} - x(n) e^{-\frac{1}{2\pi k}}$ $k = 0, 1, \dots, N - r$
 $\chi(k) = \frac{1}{N} - \frac{1}{N} = \frac{1}{N} - \frac{1}{N} = \frac{1}{N} + \frac{1}{N} = \frac{1}{N} + \frac{1}{N} = \frac{1}{N} + \frac{1}{N} = \frac{1}{N} + \frac{1}{N} + \frac{1}{N} = \frac{1}{N} + \frac{1}{N} +$$$$$$$$

EnggTree.com This expression for the FT is a polyhow of (Lagrange) interpolation formula for Xlw) expressed in terms of the values 3 xlbs } of the polynomial at a set of equally spaced @ Relationship to the Fourier Servis w-efficients of a continuous -time signal Suppose that realt) is a continuous-time periodic signal with fundamental period Tp = 1 the signal can be expressed in a Fourier Beries ralt) = & Cke jækt Fo k=-o If ralt) is sampled at a uniform rate $F_s = \frac{N}{T_p} = \frac{1}{T}$ thus $x(h) \equiv x_a(hT) = \frac{3}{2\pi} c_k e^{j2\pi k F_{0}T}$ $= \frac{5}{k_{z}-D} C_{k} e^{j\frac{2\pi k_{z}}{N}}$ $= \underbrace{\underbrace{\underbrace{\underbrace{5}}}_{k=0}^{\chi_{l-1}} \left[\underbrace{\underbrace{5}}_{l=-0}^{\chi_{0}} \underbrace{\underbrace{5}}_{k-lN} \underbrace{\underbrace{1}}_{N} \underbrace{\underbrace{5}}_{l=-0}^{\chi_{0}} \underbrace{\underbrace{5}}_{N} \underbrace{\underbrace{5}}_{N} \underbrace{\underbrace{5}}_{N} \underbrace{\underbrace{5}}_{N} \underbrace{\underbrace{5}}_{N} \underbrace{1}_{N} \underbrace{5}_{N} \underbrace{5}_$ a in the form of IDFT where egz @ $\chi(k) = N \leq C_{k-lN} \equiv NC_{k}$ $d = -\infty$ $Ch = \sum_{l=-0}^{\infty} C_{k-lN}$ Eng sequence is a aliased version of the sequence

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Properties OF DFT TI-(H44-440)
The properties OF DFT TI-(H44-440)
The properties of DFT are,
() performediate of DFT are,
() performediate of DFT are,
() charavity
() charavity
() charavity properties of DFTs and circular
convolutions
() Multiplication of two DFTs and circular
convolutions
() Circular frequency shift
() Complex conjugate properties
() Circular frequency shift
() Complex conjugate properties
() Complex conjugate properties
() Multiplication of two sequences
() Periodicity:
The rew and x(k) are an N-point DFT
point, this
$$\chi(n+N) = \chi(N)$$
 the
 $\chi(k+N) = \chi(N)$ the
 $\chi(k) = \frac{1}{2\pi N}$ the
 $\chi(k)$ the $\chi(k)$ the $\chi(k)$ the
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 $\chi(k)$ the $\chi(k)$ the $\chi(k)$ the
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(3) - Linearity.
If
$$x_{1}(w) \leftarrow \frac{DFT}{N} \times_{1}(k)$$

and $x_{n}(w) \leftarrow \frac{DFT}{N} \times_{1}(k)$
then for any nead - valued (a) complex - valued
constant a, and a_{2} ,
a, $x_{1}(w) + a_{2}t_{2}(w) \leftarrow \frac{DFT}{N} a_{1}X_{1}(k) + a_{2}X_{1}(k)$
(3) Chroular Symmetries of a sequence
The N-point DFT of a finite duration
sequence $x(w)$, of digits L2N, is equivalent
to the N-point DFT of a finite duration
 $x_{p}(w)$, of period N, which is obtained by
periodic celly extending $x(w)$,
L $x_{p}(w) = \stackrel{<}{\leq} x(w - LN)$
to here the types Another periodic requires
 $k + ke x_{p}(w)$ sequence $x + shifted$ by k with
to the tight, thes another periodic requires
 $k + with duration requires$
 $k + with duration requires
 $x'(w) = \int_{0}^{\infty} \frac{x_{1}(w)}{w} + \frac{x_{1}(w)}{w} = \frac{x_{1}(w)}{w} + \frac{x_{1}(w)}{w} = \frac{x_{1}(w)}{w} + \frac{x_{1}(w)}{w} = \frac{x_{1}(w)}{w} + \frac{x_{1}(w)}{$$

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C) Circular shift
$$= \lim_{k \to \infty} \lim_{k$$

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1) Real - Valued sequence
$$\chi(n)$$
 is real, then
 $\chi(N-k) = \chi^{A}(k) = \chi(k)$
 $|\chi(N-k)| = |\chi(k)|$ and $\mathcal{K}(N-k) = -\mathcal{K}(k)$
 $|\chi(N-k)| = |\chi(k)|$ and $\mathcal{K}(N-k) = -\mathcal{K}(k)$
(c) Real and even requescal-
If $\chi(n) \in \chi_{0}(n)$ and $u(n, n)$
 $\hat{\mu}, \quad \chi(n) = \chi(N-n), \quad 0 \leq n \leq N-1$
then $\chi_{I}(k) = 0$
 $\therefore \quad \chi(k) = \sum_{n=0}^{N-1} \chi(n) \cos \frac{2\pi kn}{N}, \quad 0 \leq k \leq N+1$
 $\psi(k) = k \quad (k + 1) \leq 0$
 $\chi(n) = \frac{1}{N} + \sum_{k=0}^{N-1} \chi(k) \cos \frac{2\pi kn}{N} = 0 \leq n \leq N+1$
then $\chi_{I}(k) = 0$
 $\chi(n) = \frac{1}{N} + \sum_{k=0}^{N-1} \chi(k) \cos \frac{2\pi kn}{N} = 0 \leq n \leq N+1$
then $\chi_{R}(k) = 0$.
If $\chi(n) \leq n \leq n \leq N-1$
then $\chi_{R}(k) = 0$.
 $\chi(k) = -\int_{n=0}^{N-1} \chi(n) \sin \frac{2\pi kn}{N}, \quad 0 \leq k \leq N+1$
 $\psi(k) \leq 0$ purely imaginally and odd.
Since $\chi_{R}(k) = 0$,
the IDFT E
 $\chi(n) = \int_{N-1}^{N-1} \chi(k) \sin \frac{2\pi kn}{N} + 0 \leq n \leq N-1$

Engineer in adjunction
If x(D & purely in adjunction
If x(D & purely in adjunction,

$$\mu = \chi(D) = \int x_T(D)$$

 $\mu = \chi_R(k) = \int_{-\infty}^{N-1} x_T(D) \sin \frac{2\pi k D}{N}$
 $\chi_T(k) = \int_{-\infty}^{N-1} x_T(D) \cos \frac{2\pi k D}{N}$
 $\chi_T(k) = \int_{-\infty}^{\infty} x_T(D) \cos \frac{2\pi k D}{N}$
 $\chi_L(k) = 0 \, dd \quad and \quad x_T(k) = 0$
 $= \chi(k) = 0 \, dd \quad and \quad \chi_T(k) = 0$
 $= \chi(k) = \mu purely \ read$
If $\chi_T(D) = \omega end , \quad then \quad \chi_R(k) = 0$
 $\Rightarrow \chi(k) = \mu purely \ read$
If $\chi_T(D) = \omega end , \quad then \quad \chi_R(k) = 0$
 $\Rightarrow \chi(k) = \mu purely \ read$
 $\chi(k) = \chi_R(D) + \chi_R(D) + \int \chi_T(D) + \int \chi_T(D) + \int \chi_T(D) + \int \chi_T(D) + \chi_R(D) + \chi_R(D) + \int \chi_T(D) + \chi_R(D) + \chi_R(D) + \int \chi_T(D) + \chi_R(D) + \chi_R(D$

5). Multiplication of Two DFTs and Circular
Convolution
Suppose that there are two finite duration
sequences of dright N,
$$x_1(n)$$
 and $x_1(n)$, this
respective $N - point DFTs$ are
 $x_1(k) = \frac{N^{-1}}{2\pi} x_1(n) = \frac{1}{2\pi N} k=0, 1, \dots, N^{-1}$
 $x_2(k) = \frac{N^{-1}}{2\pi} x_1(n) = \frac{1}{2\pi N} k=0, 1, \dots, N^{-1}$
 $x_1(k) = \frac{N^{-1}}{2\pi} x_1(n) = \frac{1}{2\pi N} k=0, 1, \dots, N^{-1}$
If $x_1(k)$ and $x_1(k)$ are multiplied then,
 $x_3(k) = x_1(k) x_1(k) = k=0, \dots, N^{-1}$
The $JDFT = 0$ f $x_3(k)$ f
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_3(k) = \frac{1}{2\pi N}$
by substituting 0 and 0 ,
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(k) x_1(k) = \frac{1}{2\pi N}$
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(n) = \frac{1}{2\pi N}$
by substituting 0 and 0 ,
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} [\frac{N^{-1}}{2\pi N} x_1(k) = \frac{1}{2\pi N} e^{\frac{1}{2\pi N}}$
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(n) = \frac{1}{2\pi N} e^{\frac{1}{2\pi N}}$
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(n) = \frac{1}{2\pi N} e^{\frac{1}{2\pi N}}$
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(n) = \frac{1}{2\pi N} e^{\frac{1}{2\pi N}}$
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(n) = \frac{1}{2\pi N} e^{\frac{1}{2\pi N}}$
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(n) = \frac{1}{2\pi N} e^{\frac{1}{2\pi N}}$
 $x_3(m) = \frac{1}{N} \sum_{k=0}^{N^{-1}} x_1(n) = \frac{1}{2\pi N} e^{\frac{1}{2\pi N}}$

EnggTree.com a=1 when m-n-l & a multiple of Al. and a =1 for any value of a = to. $\frac{N-1}{\sum} a^{k} = \int N, \quad d = m-n+pN = ((m-n))_{N},$ $k = 0 \qquad p \text{ an inleger}$ p an inlegersubstituting this in app, N-1 $p_{s}(m) = \sum_{n=0}^{\infty} x_{1}(n) x_{1}((m-n))_{N}, m = 0, 1... N = 0$ The a in the ferm of a convolution sum. The convolution sum involves the index ((m-n)) and & called circular convolution. Thus, The maltiplication of the DFTS of two sequences & equivalent to the circular convolution of the two sequences in the time domain. If x, (n) DFT x, (k) and x2(m(DFT, x2(k)) then, x, (n) (N) x, (n) (DFT, X, (k) X2 (k) than X((b-2)) a the X(b) e Jan 44

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(b) Thue represent of a require
$$r = \frac{1}{N} = \frac{1}{N} + \frac{1}{N}$$

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$$= \sum_{n \geq 0} \chi((n-d))_{N} e^{-\int_{n}^{n+n} \frac{1}{N}}$$

$$= \sum_{n \geq 0} \chi((n-d))_{N} e^{-\int_{n}^{n+n} \frac{1}{N}}$$

$$= \sum_{n \geq 0} \chi((n-d)) = \chi((n-d)) e^{-\int_{n=0}^{n+n} \frac{1}{N}}$$

$$= \sum_{n \geq 0} \chi((n-d))_{N} e^{-\int_{n=0}^{n+n} \frac{1}{N}} = \frac{d-1}{N} \chi((n-d)) e^{-\int_{n=0}^{n+n} \frac{1}{N}}$$

$$= \sum_{n \geq 0} \chi((n-d)) e^{-\int_{n}^{n+n} \frac{1}{N}} = \sum_{n \geq 0} \chi((n-d)) e^{-\int_{n}^{n+n} \frac{1}{N}}$$

$$= \sum_{n \geq 0} \chi(n) e^{-\int_{n}^{n+n} \frac{1}{N}} \chi(n) e^{-\int_{n}^{n+n} \frac{1}{N}}$$

$$= \sum_{n \geq 0} \chi(n) e^{-\int_{n}^{n+n} \frac{1}{N}} \chi(n) e^{-\int_{n}^{n$$

9). Complex - Configure properties:

$$\chi(n) < \frac{DF}{N} × (k)$$

$$+ \lim_{N \to \infty} x^{+}(k) = \sum_{n} ((-k))_{N} = x^{+}(n-k)$$

$$The FDFT = q + x^{+}(k) = \sum_{n} (-\sum_{n=1}^{N+1} x^{+}(k) = \sum_{n=1}^{n} (-\sum_{n=1}^{N+1} x^{+}(k) = \sum_{n=1}^{N+1} x^{+}(k) = \sum_{n=1}^{N+1} x^{+}(k) = \sum_{n=1}^{N+1} x^{+}(k)$$

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By evaluting IDFT at 1=0, the equip @ can be obtained. The equin (D) is called the general form of the por seval's theorem. Jathan y (m= x (m) then N-1 $\frac{5}{120} |x(n)|^2 = \frac{1}{11} \frac{5}{5} |x(k)|^2$ which is the energy in the finite - duration sequence kins in terms of the foreguency components fx (b) } 4) Find the linear and circular convolution of the two sequence x, (n) = { 1,2,2,2 } and ×2107 = [1,2,3,4]. (AIM-17- ROF) Linear convolution: x, (2) 3 (1) 14 2 - (1) 2 1 3 2 x1 m 8 6 4 2 6/08 001 LP 2 y(n) = \$ 1, 4, 9, 16, 18, 14, 87 reg (L) = to 5 - kay (L) er



5

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$$y_{1}(k) = \int_{1}^{2} \tau_{1} - 1 \cdot -1 \int_{1}^{2} \frac{1}{1} = 0$$

$$y_{2}(k) = \int_{1}^{2} \tau_{2}(k) = \int_{1}^{2} \frac{\pi k h}{h} = 0$$

$$y_{2}(k) = \frac{3}{h_{2}} + \frac{1}{h_{2}} + \frac{1}{h_{2}} = 10$$

$$y_{2}(k) = 1 + 2 + 3 + 4 = 10$$

$$y_{2}(k) = 1 + 2 + 3 + 4 = 10$$

$$y_{2}(k) = 1 + 2 + 3 + 4 = 10$$

$$y_{2}(k) = 1 + 2 + 3 + 4 = 10$$

$$y_{2}(k) = 1 + 2 + 3 + 4 = 10$$

$$y_{2}(k) = -2 + 3 - 4 + \frac{1}{h_{2}} + \frac{$$

9

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$$x_{3}(i) = \frac{1}{h} \sum_{k=0}^{3} x_{4}(k) = \int_{1}^{1} \frac{1}{h^{2}} \sum_{k=0}^{3} x_{4}(k) = \int_{1}^{1} \frac{1}{h^{2}} \sum_{k=0}^{3} \sum_$$

$$\begin{array}{l}
 B & b) & The first first DFT co-afficients of a sequence R(D) are $X(D) = 20$, $X(D) = 5+j^{2}$
 $X(D) = 0$, $X(D) = 0.2+j^{0} + X(D) = 0$.
Determine the remaining DFT coefficients.
 $(H/J - 0T - Rod)$
Set:
 $X(K) = \int 20, 5+j^{2}, 0, 0.2+j^{0.4}, 0, K(D), X(D)$
 $X(K) = \int 20, 5+j^{2}, 0, 0.2+j^{0.4}, 0, K(D), X(D)$
From the complex conjugat property of DFT.
 $X(K) = X^{*}(N-K)$.
 $(T) = X^{*}(B-5) = X^{*}(D) = 5-j^{2}$
 $X(K) = \int 20, 5+j^{2}, 0, 0.2+j^{0.4}, 0, 0.2-j^{0.4}, 0, 0$
 $X(T) = X^{*}(F-T) = X^{*}(D) = 5-j^{2}$
 $(X(K) = \int 20, 5+j^{2}, 0, 0.2+j^{0.4}, 0, 0.2-j^{0.4}, 0, 0)$
 $(T) = (X^{*}(F-T) = X^{*}(D) = 5-j^{2})$
 $(X(K) = \int 1, 0, 1, -1, 1, -1, 0, 1 \int (N/D-12, Kos) + (1-1)5-j^{2} \int (N/D-12, Kos) + (N/D-12, Kos) + (1-1)5-j^{2} \int (N/D-$$$

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FILTERING METHODS BASED OF DET

Since DFT provides a déscrete frequency representations of a finite - duration requence in the frequency domain, it is used as a computational tool for linear system analysis and especially for linear filturing.

Use of DFT in linear filtering:

3

The product of two DFTs & equivalent to the circular convolution of the corresponding true-domain sequences. Unfortunally, circular convolution & of no use, if the objective & to find the output of a linear falter to a given input sequence. In this case a frequency domain methodology equivalent to linear for convolution & required.

Let XIND be the finde devoltor sequence XVAD of dength L which excites on FIE filler of length M. Let XIND = 0 nCO & N > L

where him is the impulse response of the filter.

anthora-

The output sequence;

$$g(n) = \underset{k=0}{\overset{M-1}{\underset{k=0}{\overset{}}} h(k) \times (n-k)$$

The length of yours and had and get

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In frequency domain,

$$71w = x(w) + 1(w)$$

If $91(k) = 91(w)/w = 2\pi k$, $k = 0, 1, ..., n = 1$
 $= x(w) + 1(w) / w = 2\pi k$, $k = 0, 1, ..., n = 1$
 $= x(w) + 1(w) / w = 2\pi k$, $k = 0, 1, ..., n = 1$
 $= x(w) + 1(w) = x(k) + 1(k)$, $k = 0, 1, ..., n = 1$
then $9(k) = x(k) + 1(k)$, $k = 0, 1, ..., n = 1$
 $4w + 1(k) = x(k) + 1(k)$, $k = 0, 1, ..., n = 1$
 $4w + 1(k) = x(k) + 1(k)$, $k = 0, 1, ..., n = 1$
 $4w + 1(k) = x(k) + 1(k)$, $k = 0, 1, ..., n = 1$
 $4w + 1(k) = x(k) + 1(k)$, $k = 0, 1(k)$, $h(k)$ and
 $4w + 1(k) = x(k) + 1(k)$, $2w = padd by can be done have.$
Thus with: zero padd by DFT can be used
 $4w + 1(k) = 2w = padd by DFT can be used
 $4w + 1(k) = 2w = padd by DFT can be used
 $4w = 1(k) = 2w + 1(k) + 1(k$$$
EnggTree.com T 1 0 0 2 -2 1+0+0+0+0=1 2 0 2 0+0=0 A 1 0 0 2 -2 12 0 0 2 Loo 0 0 + 0 + 2 + 0 + 0 = + 2y(n)={ 1, 0, -1, 2, +2} To verify_using dinear convolution him lapient - 2" 2 promo m there if i send twith larges liqui gral A Aire by blocks prior of 2102 fillening & Direar, our course 1 0, -1, 2, 24 タレのニティ, into black the reput data requere a of L paints where L>>H # The size of the input blacks . & N=L+M-1 and DFTS and IDFT are Each date block consuls of eltre last Mai deta points of the memory data black

FILTERING OF LONG DATA SEQUENCES [M/J-4 Ros] In practical applications involving linear filtering of signals, the input sequence KIN & signed monistoring and analysis. Since linear filtering performed Via the

DET involves operations on a block of black, which is by necessity must be dimited in size due to dimited memory of a digital computer. A doing input signal must be segmented to if is and size by blocks prior to processing. Since the fidturing is dimeer, successive blocks can be processed one at a time via the DET, and the output blocks are fitted together to form the overall output signal sequence.

Overley-save and Overlep-add methods are used for long data sequences.

Apsume FIR filler bes devalues M, The input data sequence to segmented into block of L points where L>>M.

1) OVERLAP - SAVE METHOD! # The size of the input blocks. & N = L + M - 1 and DFTs and IDFT. are of length N. # Each data block 'constants of the last M-1 data points of the previous date block

Followed by L new dave prime to first to first a
data sequence of length N= L+H-1.
* N-point DFT & computed for each data block
* The impulse response of the FIR filter a
intreased in length by apponding L-1 zeros
and N-point DFT & computed once and stored.
* The multiplication of the two N-point DFTs

$$j + (W) j$$
 and $j X_m(W) j$ for the mith block
of data yield
 $f_m(W) = H(W) X_m(K)$ $k=0,1...,N-1$
Then the N-point IDFT fields the routh.
 $j'_m(W) = j Y_m(W) j$ for the next block
 $d'_m(W) = H(W) X_m(K)$ $k=0,1...,N-1$
Then the N-point IDFT fields the routh.
 $j'_m(W-1) j$
* Since the data record is of dargte N.
Its first H-1 points of Jm(N) are corrupted
by aliasing and must be discarded.
* The least L point of Jm(N) are excleding
its same as the routh from dimar convolution
 $j'_m(W) = j'_m(N)$, $n = M$, $H+1$... $N-1$
* To avoide loss of data due to alicency.
its lost H-1 points of condents are saved
and these points before the first M-1 data points
of discongenust record.
* To short the proof, the first M-1 data points
of discongenust record.

(R)

F. F.

Thus the blocks of data sequences are

$$x_1(n) = \begin{cases} 0, 0, \dots 0, 2(0), 2(1) \dots 2(L-1) \end{cases}$$

 $M - 1 points$

$$\chi_2(n) = \begin{cases} \chi(L-H+1), \dots \chi(L-1), \chi(L), \dots \chi(2L-1) \end{cases}$$

$$M - 1 data points from L new data points
\chi_1(n)$$



The two XI-points DFTs are multiplied togethe to form $Y_m(h) = H(h) X_m(h)$ k=0,1...,N-1The IDFT yields data blocks of length N that one free of aliasing, since the size of the DFTS and IDFT. & N=L+M-1 and the sequences are increased to Ni-petito by appending zeros to each block. of since each data block to terminated with

M-1 zeros, the lest M-1 points from each output block must be overlapped and added to the first M-1 points of the succeeding block.

$$\begin{split} \dot{u}, & y(n) = \begin{cases} y_1(n), y_1(n), \dots, y_1(n-1), y_1(n) + y_2(n) \\ & y_1(n+1) + y_2(n), \dots, y_1(n-1) + y_2(n-1) \\ & + y_2(n), \dots, y_1. \end{split}$$

9) Perform linear convolution of finite durates sequences $h(n) = \frac{1}{2} 1, 2\frac{3}{2}$ and $x(n) = \frac{1}{2} 1, 2, -1, 2, 3, -2, -3, -1, 1, 1, 2, -\frac{1}{2}$ by overdap add and save method.

Overdap save Method

L771)

(NID-16 R#3) (NID-11 Ros)

Let
$$f = 3$$
 Engg[Jee.com_{-1} = 4 the data block
Divide the head sequence $x(n)$ but data block
as shown below.
 $x_{1}(n) = \int_{-1}^{n} 0, \quad 1, 2, -1, \frac{1}{2}$
 $H_{1} = f + \frac{1}{3}$ new data
 $x_{2}(n) = \int_{-1}^{n} -1, 2, \frac{1}{3}, -2, \frac{1}{3}$
 $H_{1} = f + \frac{1}{3}$ new data
 $data from
previous block
 $H/40; \quad x_{3}(n) = \int_{-2}^{n} -2, -2, -1, \frac{1}{3}$
 $2\mu(n) = \int_{-1}^{n} 1, \frac{1}{2}, \frac{2}{-1}, \frac{1}{3}$
 $x_{1}(n) = \int_{-1}^{n} 1, \frac{1}{2}, \frac{2}{-1}, \frac{1}{3}$
 $x_{2}(n) = \int_{-2}^{n} 1, \frac{1}{2}, \frac{2}{-1}, \frac{1}{3}$
 $x_{2}(n) = \int_{-2}^{n} 1, \frac{1}{2}, \frac{2}{-1}, \frac{1}{3}$
 $x_{2}(n) = \int_{-2}^{n} 1, \frac{1}{2}, \frac{2}{-1}, \frac{1}{2}$
 $x_{2}(n) = \int_{-2}^{n} 1, \frac{1}{2}, \frac{2}{-1}, \frac{1}{2}$
 $y_{1}(n) = x_{1}(n) (n) (n) (n) (n)$
 $\int_{-1}^{n} \frac{1}{2} \int_{-1}^{n} \frac{1}{2} \int_{-1}$$

(43)

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$$y_{3}(n) = x_{3}(n) \bigoplus h(n)$$

$$\begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 & 1 & 0 & 0 \\ 0 & 2 & 1 & 0 \end{bmatrix} \begin{bmatrix} -2 \\ -3 \\ -1 \\ 1 \end{bmatrix} = \begin{bmatrix} -2 + 0 + 0 + 1 \pm 0 \\ -4 - 3 + 0 + 0 \pm -7 \\ 0 - 6 - 1 + 0 \pm -7 \\ 0 + 0 - 2 \pm 1 \pm -1 \end{bmatrix}$$

$$y_{3}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 & 1 & 0 & 0 \\ 0 & 2 & 1 & 0 \end{bmatrix} \begin{bmatrix} 1 + 4 \\ 1 \\ 2 \\ 0 & 0 & 2 & 1 \end{bmatrix} = \begin{bmatrix} 1 + 0 + 0 - 2 \pm -1 \\ 2 + 1 + 0 + 0 \pm 3 \\ 0 + 0 + 1 \pm -1 \pm 3 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 & 1 & 0 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} -1 \\ 2 \\ -1 \end{bmatrix} = \begin{bmatrix} 1 + 0 + 0 - 2 \pm -1 \\ 2 + 1 + 0 + 0 \pm 3 \\ 0 + 0 + 1 - 1 \pm 3 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 & 1 & 0 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} -1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -1 \\ -2 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 \pm 1 & 0 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} -1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -1 \\ -2 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 \pm 1 & 0 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} -1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -1 \\ -2 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 \pm 1 & 0 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} -1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -1 \\ -2 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 \pm 1 & 0 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -1 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 & 0 & 0 & 2 \\ 2 \pm 1 & 0 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} -1 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{2}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{2}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{2}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{1}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{2}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{2}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$y_{2}(n) = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

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Overlap add HetEnggTree.com

$$M = 2 \qquad Let L = 3.$$

$$M = 4$$

$$The book scepture can be duided let date
blocks as obtain below.
$$u_{1}(x) = \int_{0}^{x} \frac{1}{2} (x, -1) \frac{d}{2} \frac{d}{3} \frac{d}{3} \frac{d}{3} \frac{d}{4} \frac$$$$

1

EnggTree.com For each value of a direct computation of X(k) involves N complex multiplications 1 4N real multiplications] (N-1) complex addition [4N-2 real additions]. ... To compute N-point DFT, N² complex multiplications and N(N-1) complex additions required. () Divide - and - conquer approved (DFT of length w). Twiddle Factor The complex number $W_N = e^{-\int \frac{2\pi}{N}}$ Loy $W_N = 1 \frac{1 - 2\pi/N}{2}$ is called the phase factor (er) twiddle factor. It represents Nth root of unity. Properties of hitddle factor are, Symmetry Property: $W_N^{k+\frac{N}{2}} = -W_N^{k}$ Periodicity Property: WN = WN compate day sequence { x(b) } of N complex - value * FFT-Fast Fourier Transform to alter experient way of calculating DFT. XILD = S xLAI WO

Bradder -> FFT AFTEGGIISERCOM
The drugth N ay DFT is represented
as N = t when r> tadd x of DFT.
radie -> algorithms are nost widdy used.
I Reduct N = 2¹⁰
X Split the N-point date segume into two
NL point date sequence
$$\psi_1(n)$$
 and $\psi_2(n)$.
involved the N-point date sequence into two
NL point date sequence $\psi_1(n)$ and $\psi_2(n)$.
involved the sequence $\psi_1(n)$ and $\psi_2(n)$.
If $(n) = x(2n)$
 $\int_{1}(n) = x(2n)$
 $\int_{1}(n) = x(2n)$
 $\int_{1}(n) = x(2n)$
 $\int_{2}(n) = \frac{1}{n = 0} x(n) W_N^{h} + \sum_{n \ge 0} x(n) W_N^{h}$
 $WkT W_N^{h} = W_{NL}$
 $\int_{2}(n) = x(2n) W_{NL}^{h} + W_{NL}^{h} = x(2n+1) W_N^{h}$
 $= F_1(k) + W_N^{h} F_2(k)$, $k = 0, \dots, N-1$

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where
$$F_{1}(k)$$
 and $F_{2}(k)$ are N_{2} point DFTs
af the sequences $f_{1}(n)$ and $f_{2}(n)$ respectively.
 $F_{1}(k) = F_{1}(k)$ and $F_{2}(k + \frac{M}{2}) = F_{2}(k)$
also, $W_{N}^{k+\frac{M}{2}} = -W_{N}^{k}$
 $F_{1}(k + \frac{M}{2}) = F_{1}(k)$ and $F_{2}(k + \frac{M}{2}) = F_{2}(k)$
also, $W_{N}^{k+\frac{M}{2}} = -W_{N}^{k}$
 $F_{1}(k) = F_{1}(k) + W_{N}^{k}F_{2}(k)$
 $k = 0, 1, \dots, \frac{M}{2} - 1$
 $K(k + \frac{M}{2}) = F_{1}(k) - W_{N}^{k}F_{2}(k)$
 $k = 0, 1, \dots, \frac{M}{2} - 1$
The divise computation of $F_{1}(k)$ and $F_{2}(k)$
 $requires$ $(\frac{M}{2})^{2}$ complete multiplications complet
multiplications required its computer $W_{N}^{k}F_{2}(k)$
 $F_{2}(k) + M_{2} = \frac{N^{2}}{2} + \frac{M}{2}$
complex multiplications.
 $F_{2}(k) + \frac{M}{2} = \frac{N^{2}}{2} + \frac{M}{2}$
complex multiplications.
 $F_{2}(k) + \frac{M}{2} = f_{2}(k) - K + \frac{M}{2}$, which is about
a factor of 2 for N longe.
Let $C_{1}(k) = F_{1}(k)$ $k = 0, 1 \dots \frac{M}{2} - 1$
Thus, $\chi(k) = C_{1}(k) - C_{2}(k)$

Having performed its decimation - in-time once,
the process can be repeated for each of the
sequence
$$f_1(\omega)$$
 and $f_1(\omega)$. Thus $f_1(\omega)$ works result
in the two $\frac{M}{4}$ point sequence.
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $u_{in}(n) = f_i(2n)$ $n = 0, 1 + \frac{M}{4} - 1$
 $H_{in}(n) = f_{in}(n)$ $h = 0, 1 + \frac{M}{4} - 1$
 $h = 0, 1 + \frac{M}{4} - 1$
By computer $f_{in}(n) = \frac{M}{4}$ point $BFTS$, the $\frac{M}{4}$ point
 $F_{in}(k) = V_{in}(k) - W_{M/k}(k)$, $k = 0, 1 + \frac{M}{4} - 1$
 $F_{in}(k) = V_{in}(k) - W_{M/k}(k)$, $k = 0, 1 + \frac{M}{4} - 1$
 $F_{in}(k) = V_{in}(k) - W_{M/k}(k)$, $k = 0, 1 + \frac{M}{4} - 1$
 $F_{in}(k) = V_{in}(k) - W_{M/k}(k)$, $k = 0, 1 + \frac{M}{4} - 1$
 $F_{in}(k) = V_{in}(k) - W_{M/k}(k)$, $V_{in}(k)$
 $F_{in}(k) = V_{in}(k) - W_{M/k}(k)$, $V_{in}(k)$
 $F_{in}(k)$ requires
 $\frac{M_{in}^{2} + M_{in}}{k} + \frac{M_{in}}{k} = \frac{M_{in}^{2}}{k}$.
The decimation of k date asymme can be
repeated spin and spin with star rawards requires
are reduced to on point sequence.

* For N=2" this decimation can be performed Le = log N times. * Thus the total number of complex multiplication regioned & reduced to $\frac{N}{2}$ dog N Block diagram of 8-point DIT-FFT



For N=8, the computation is performed in three stages, beginning with computation of your 2-point DFTs, this two goss 4-point DFTs and finally one 8-point DFT.

Basic Butterfuly diagram of DIT-FFT



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(k) EnggTree.com 8-point DIT-FFT Algorithm:x (0) " X(0) WX xly . X(1) Wg x12) . X(2) W8 x(6) XQ) W80 (2) x (4) Wr WE X(5) WE WS xB) X(6) Wr Wr x(7) . 5 X(7) Jacobian position dem Gas join jim Each Butterfuy involves one complex multiplication and ition compiler addictions. For N= 2 there are N/2 butterflies per stage of the compution process and log N stages. ". The total number of complex multipliceetions required & (N/2) log N The dotal number of complex additions required in NI log 2 00----

In -place computation

2N storage the registers are required to store the results (N complex numbers) of the FFT computations at each stage. Since the same 2N storage locations are used throughout the computation of the N-point DFT, the Computation & named as in-place computation.

Bit - Reversal :

XED

The order of the decimated date sequence. in FFT & obtained by reading the binary representation of the index n in reversed order. Ex, the deta point $r(3) = r((011)_2) &$ placed in position m= (110), con m=6 in the decimated array. ... The data relad offer decimation & stored in bit-reversed order. This is called bit reversal.



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$$R_{BDIS} = 2 \quad D \mid F - FFT \quad Alequelitum \quad [m/5 - 15 \ kor]$$

$$D \mid F - D evimation \quad In \quad Fraquency$$

$$splid the DFT france into two summations
of which one involves the sum over the dist N/s
dele point and the second, sum over the dest
$$N/s \quad dele point.$$

$$Thue, \quad \frac{M-1}{N-1} \quad x(n) \quad W_N^{h} + \frac{N-1}{2} \quad x(n) \quad W_N^{h}$$

$$= \frac{M}{2} \quad x(n) \quad W_N^{h} + \frac{N+1}{2} \quad x(n) \quad W_N^{h}$$

$$= \frac{M}{2} \quad x(n) \quad W_N^{h} + \frac{N+1}{2} \quad x(n + \frac{M}{2}) \quad W_N^{h}$$
Since $W_N^{hF} = (-D^{h}, \quad He expression can be written as
$$X(k) = \frac{M-1}{2} \quad \left[x(n) + (n + \frac{M}{2}) \right] W_N^{h}$$
Now split $X(k)$ into two and odd numbered
sample.

$$X(2k) = \frac{M-1}{2} \quad \left[x(n) + x(n + \frac{M}{2}) \right] \quad W_N^{h} \quad k \ge 0, \dots, \frac{M-1}{2}$$

$$\int W_N^{h} = \sum_{h=0}^{M-1} \quad \left[x(n) + x(n + \frac{M}{2}) \right] \quad W_N^{h} \quad k \ge 0, \dots, \frac{M-1}{2}$$

$$\int W_N^{h} = \sum_{h=0}^{M-1} \quad \left[x(n) + x(n + \frac{M}{2}) \right] \quad W_N^{h} \quad k \ge 0, \dots, \frac{M-1}{2}$$

$$\int W_N^{h} = \sum_{h=0}^{M-1} \quad \left[x(n) + x(n + \frac{M}{2}) \right] \quad W_N^{h} \quad k \ge 0, \dots, \frac{M-1}{2}$$

$$\int W_N^{h} = \sum_{h=0}^{M-1} \quad \left[x(n) + x(n + \frac{M}{2}) \right] \quad W_N^{h} \quad k \ge 0, \dots, \frac{M-1}{2}$$

$$\int W_N^{h} = \sum_{h=0}^{M-1} \quad \left[x(n) + x(n + \frac{M}{2}) \right] \quad W_N^{h} \quad k \ge 0, \dots, \frac{M-1}{2}$$

$$\int W_N^{h} = W_{h}$$

$$ff \quad g_1(n) \quad and \quad g_2(n) \quad one \quad He \quad two \quad M \quad point sequence,$$

$$g_1(n) = x(n) + x(n + \frac{M}{2})$$$$$





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Productive Obstrag DIT-FET ALGORITHM
10) Find DFT of the import acquire
$$r(n) = \int 0, 1, s, d$$

why DIT-FET algorithm.
See Given $N = 4$ $r(n) = \int 0, 1/s, s^{2}$
 $W_{N}^{0} = s$ $W_{N}^{0} = s^{-1}$
 $W_{N}^{0} = s$ $W_{N}^{0} = s^{-1}$
 $r(0) = 0$ $x(0) = 6$
 $r(0) = 0$ $x(0) = -2stys$
 $r(0) = 0$ $r(0) = -2stys$
 $r(0) = -2stys$



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EnggTree.com PROBLEMS COUNTR DNE-FET ADGODITHING x-1 12) Find DFT of the sequence xlore for 1, -1, 1, -1, using DIT-FFT algorithm. Given, $x(n) = \int \frac{1}{67773} \int$ WN = e J2TK $W_2^0 = I$ $W_4^\dagger = -j$ 72(0)=1 -X (0)=2-2=0 x[2]_+ W4 = 1 21 -x10= 0 250=-1 W4P=1 -2 W4°=1 -1 -1 - × (2)= 4 0 W4'=-- x(s)= 0 · · | ×(k)= { 0,0,4,0} PROBLEMS OSING DIF-FFT ALGORITHM 13 Find DFT of the sequence xm) = { 1, -1, 1, -1 } using DIF-FFT algorithm. Given, 210) - \$ 1, -1, 1, -17 $\frac{k_{1}}{w_{N}} = e^{-\int \frac{2\pi k}{N}}$ $W_4^\circ = 1$ $W_{4}' = -j$





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and a start of the EnggTree.com UNIT-II IIR FILTER DESIGN Structures of JIR - Analog filter design - Aberete time IIR filler from analy filler - II & filler design by Impube Invariance, Bilinear transformation, Approximation of derivatives - (LPF, HPF, BPF, BRF) filter design using frequency translation. TI-(101-703) IIR Fidters - Infinite Impube Response Filters Digital filter 20 a linears time invariant discrete time system. FIR fillers - Non recursive type - present output sample depends on the present input sample and previous input sample. IIR filters - Receive type - present output sample depends on the present input past imput and past output also. i = y(n) = x(n) + x(n-1) + ...+ y 10-1) + y(0-2)+... Take z transform, 4(2)=box(2)+biz 1x(2)+... +az-1 Y(z)+az Y(z)+... $Y(z) - az - (Y(z) - az - (Y(z) - b_0 X(z) + z - (X(z) + z - X(z)) + z - (X(z) + (X(z) +$ $\frac{Y(z)}{x(z)} = \frac{b f_{0} z^{-1} + b z^{-1$ 1+(2)= = bk2-k -> x 1+ = ak2-k ->)

EnggTree.com They the trave of IIA follow will be. H(2) = Kop bk a 1+ 5 ak 2-4 The design of JIK falter for the given specification is finding the of the co-off. ak and bk. Design of digital filters from Analog filters ; step - Map the disired digital filter specification into those for an againalent analog filler. about Derive the analog transfer function for etter analog protopype. sitions Transform the transfer function of ette analyg protogoe into an equivalent degreat filter transfer fenctros. Magnitude Response of LPF: INCEST Digital Analy (megan) Te p-B SB 00 00 Ac sis aparameter specifying allowable pauband es 2 150 A -> parameter specify of allowable stopsand A= K+ 6p) - 65

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Analog filters
$$\int_{CLE} Bulterworth filter
Ls ris Order (W)
Ls (ii) Transfer function (Ha (D))
Analog Bulterworth filter duign: Ti (1/1-710)
(i) Order of the filter (N)
The magnitude function & gives by,
 $\left| H(\frac{1}{2}2) \right|^{2} = \frac{1}{\left[1 + e^{2} \left(\frac{-2\pi}{2}/2p\right)^{2N}\right]^{1/2}}$
 $\left| H(\frac{1}{2}2) \right|^{2} = \frac{1}{\left[1 + e^{2} \left(\frac{-2\pi}{2}/2p\right)^{2N}\right]^{1/2}}$
 $\left| H(\frac{1}{2}2) \right|^{2} = \frac{1}{(1 + e^{2} \left(\frac{-2\pi}{2}/2p\right)^{2N}}\right]^{1/2}$
Taking uslogs on both andu,
 $\left| Udg(\left[H(\frac{1}{2}2n)\right]^{2} \right] = \left[Udg(1 - \log \log \left[\ln e^{2}/2p\right]^{2N} \right]$
 $20 \log \left[H(\frac{1}{2}2) \right] = \left[Udg(1 - \log \log \left[\ln e^{2}/2p\right]^{2N} \right]$
 $h \to = p_{2}$, the attanuation is αp .
 $\therefore so dg(1 + (\frac{1}{2}2n)) = -\alpha_{2}$
 $\mu \ll p = \mu \left[0 \log \left[1 + e^{2} \right]^{2} \right]$
 $\left[0 \cdot 1 \ll p = \log (1 + e^{2}) \right]$
 $\left[0 \cdot 1 \lll p = \log (1 + e^{2}) \right]$$$

EnggTree.com At n=ns, the minimum skyband aftermation be de. : 20 log | H(Jaw) = 10 log 1 - 10 log [1+e2 (a))" × × = × 10 log (1+ e 2 (-03) 2 m) 0.1 xs = log 1+ e2 (23) 2N] $\frac{0.18}{10} = 1 + E^2 \left(\frac{-2}{2p}\right)^{2N}$ C2 (55) 2N = 10 0.123-1 $\left(\frac{-r_{s}}{-r_{p}}\right)^{n} = \frac{10^{\circ} - 1}{10^{\circ} \cdot 1 \times p}$ Take log on both sides, 0, N log (-2) = log 10 0.12 -1 10 -1 -1 10 -12 -1 $N = \frac{\log \int \frac{10^{\circ 1 \times 3} - 1}{10^{\circ 1 \times 3} - 1}}{10^{\circ 1 \times 3} - 1}$ log (- sp) The above expression doesnot result in an Inter value. $N \geq \frac{\log\left(\frac{10^{0.1K_{s}}-1}{10^{0.1K_{p}}-1}\right)}{\log\left(\frac{10^{0.1K_{p}}-1}{10^{0.1K_{p}}}\right)}$

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N Z Log &

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where
$$\lambda = \sqrt{10^{\circ 1/4} - 1}$$

 $e = \sqrt{10^{\circ 1/4} - 1}$
 $e = \sqrt{10^{\circ 1/4} - 1}$
 $h = \frac{1}{e}$ and $h = \frac{\pi}{1 + 1}$
 $h = \frac{1}{e}$ and $h = \frac{\pi}{1 + 1}$
 $h = \frac{1}{e}$ and $h = \frac{\pi}{1 + 1}$
 $h = \frac{1}{1 + 1$

This above equations fells that it has polle in the LHP (Left Half of a Pilane) as well as in the RHP (Right Half of a Pilane), because of the two factors H(s) and H(-s).

If HG) has roots in the LHP then HE-s) has the corresponding roots in the KHP.

The roots can be obtained by equating the demonsinator to zero.

ie, 1+(-s²)^N = 0. For N-odd, the above eyes reduces to

8 = e	To ensure when
$8k = e^{\int \frac{\pi k}{N}}$	k=1,2,2N

For N-even, the egg & reduces to.

$$S^{2N} = -1 = e^{\int (2k-1) \pi}$$

 $\int \frac{j(2k-1)\pi}{2N} = \int for \ k=1, 2, ... 2N$

For M=3, $3^6=1$

For
$$N = 0.000$$
,
 $S_1 = e^{-\int \frac{\pi}{3}} = 0.5 + \int 0.866$
 $S_2 = e^{-\int \frac{2\pi}{3}} = -0.5 + \int 0.866$
 $S_3 = e^{-\int \pi} = -1$
 $S_4 = e^{-\int \frac{4\pi}{3}} = -0.5 - \int 0.866$
 $S_5 = e^{\int \frac{5\pi}{3}} = 0.5 - \int 0.866$
 $S_5 = e^{\int \frac{5\pi}{3}} = 0.5 - \int 0.866$


All the poles are located in the s-plane. The angular separation between the poles is $\frac{360}{2N}$ is 60° and all the poles die on a unit circle.

To ensure stability, only the poles that lie on the left half of the spilene are considered.

... The dimonishador of the transfer feusclios

$$(3+i) \stackrel{2}{\leq} (3+i) \stackrel{2}{\leq}$$

$$H(s) = \frac{1}{(s+1)(s^2+s+1)}$$

As we are interested on the poles, which lies in the left half of the s-poleine, the same can be found by using where, $q_k = \frac{T}{2} + \frac{(2k-D)T}{k} = 1, 2, ... N$

EnggTree.com Lost of Butterworth Polynomials N Denominator of H(S) 8+1 1 S2 + J2 3 +1 2 (0+1) (02+ 0+1) 3 (32+0.765373+1) (32+1.84778+1) 4 (3+1) (32+0.618033+1) (32+1.618033+1) 5 (32+1.9318553+1) (32+55+1) (32+0.517643+1) 6 (S+1) [S2+1.801945+1) (32+102473+1) 7 (32+0.4453+D) These podes are called normalized podes, because sic = 1 rad (see. In general, its unnormalized podes are given sk'= she sh. 58 The transfer fundios of such type of Butterworth filter can be obtained by substituting 3 -> 2 In the transfer function of Balterworth filler. Frequency Transformation 1 Low pass filter to Lew pass filter 3-3-5 nosc De= D Jou pour filter to High pour filter 3 → o (Low pay filter to Band pass filter S > Arc 3(A2-A) (Low pour filter to Band stop fills S > 2 (3 + 212 +)

O Determine the order and the pole of loopas Butterworth filter that has a side attenuation at 500 the and an attenuation of Lodb at 1000 the Sop Cliven up=3db fp=500Hz as = 400B fs = 1000 Hz _p= 2×π× 500 = 1000π rad/see Q++64814. As= 2× T × 1000 = 2000 T rad Bee. Order, $x \ge \log \sqrt{\frac{1001}{1001}}$ Log <u>so</u> $\geq \log \sqrt{\frac{104}{10^{-3}}}$ Log 2000T Arcadias of and 2 6.6 N = 7 2) Prove that $\Omega_{c} = \frac{\Omega_{p}}{(10^{0.1} M_{p})^{1/2N}} = \frac{\Omega_{s}}{(10^{0.1} M_{s})^{1/2N}} \frac{1}{(10^{0.1} M_{s})^{1/2N}}$ Magnichide squeere function of Butterworth analog f.p.f. b, in a man in the $| H(jo)|^2 = \frac{1}{1+\left(\frac{n}{-2c}\right)^{2N}}$ D Wee St.



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EnggTree.com TI (719-728) Chebyshev Filters Two styges of chebyshes filters All pode filture that exhibit Type I aquiripple behavior in the pairbany and monotonic characteristics in the stop band. Type - I -. Contains both poles and Europ and exhibits a monotonic beheator in the pass band and an equipple beheator in the stepband. Zeros lie on the imaginary axts. 1Har) HUDDA Type I ap so a 0-) Type -J The magnitude squeere function of with order Type-I filter can be expressed as, 1492)1= -1 N= 1, 2, ... $1 + \epsilon^2 c_N^2 \left(\frac{a}{c_P}\right)$ where _> parameter related to the ripple in the passband CNW > Nth order Chetyshew palynomial

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$$C_N(x) = \int cos \left[N cos^{-1}x \right] \quad |x| < 1 \rightarrow possbond$$

 $\int cos h(N cosh^{-1}x) \quad |x| > 1 \rightarrow stopband$
 $C_N(x) = 2x C_{N-1}(x) - C_{N-2}(x) \quad N > 1$
 $C_0(x) = 1 \quad C_1(x) = x$

Propulties of Children pottromate.
5
$$C_{N}(x) = -C_{N}(-x)$$
, Nodd
 $C_{N}(x) = C_{N}(-x)$, Neven
 $C_{N}(x) = C_{N}(-x)$, Neven
 $C_{N}(x) = C_{N}(-x)$, Neven
 $C_{N}(x) = (-1)^{M/2}$, Neven
 $C_{N}(x) = 0$, Nodd
 $C_{N}(x) = 1$, Neven
 $C_{N}(x) = 1$, Neven
 $C_{N}(x) = 1$, Neven
 T $C_{N}(x) = -1$, Nodd
 T $C_{N}(x)$ Oscillates with equal rippule believen ± 1
for $1x1 \leq 1$
 T $C_{N}(x)$ Oscillates with equal rippule believen ± 1
 $for $1x1 \leq 1$
 T $C_{N}(x)$ $C = 1 c_{N}(x) 1 \leq 1$ for $0 \leq 1/21 \leq 1$
 $for $1x1 \geq 1$
 T $C_{N}(x) \in 0$ monotonically measured for $1x1 > 1$ for
 $all N$.$$



The pales of the chebyshee transfer function are located on an ellipse in the s-plane. The equation of the ellipse to given by,

a, 5 one ninor and major and of the ellipse.



Type-IL

It has both pales and zeros.

$$|H(j_{\mathcal{P}})|^{2} = \frac{1}{1+\epsilon^{2}} \left[\frac{c_{N}\left(\frac{n_{0}}{n_{p}}\right)}{c_{N}\left(\frac{n_{0}}{n_{p}}\right)} \right]$$

The zeros are located on the imaginery axes at the points,

Pales are located at the points (xk, yk)



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(3) Criven the specification
$$a_{p} = 3db$$
, $a_{u} = bdb$
 $f_{p} = 1kH_{2}$ and $f_{s} = 2KH_{2}$. Determine the
order of the chebysher fulter.
Solver:
 $a_{v} = 3db$ $f_{p} = 1KH_{2}$
 $a_{v} = bdb$
 $f_{p} = 1KH_{2}$
 $f_{p} = 1KH_{2}$
 $f_{p} = 1KH_{2}$
 $a_{v} = bdb$
 $a_{v} = bdb$
 $a_{v} = bdb$
 $a_{v} = bdb$
 $a_{v} = 3db$
 $a_{v} = 3db$

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$$N \ge \frac{\cos h^{-1} \frac{h}{e}}{\cos h^{+1} \frac{h}{h}} \qquad A \ge \frac{\cosh h^{-1} \left[\int \frac{1}{10} \frac{1}{10} \frac{h^{-1}}{10} \right]}{\cosh h^{-1} \left(\frac{1}{10} \frac{h}{10} \right)}$$

$$N \ge \frac{1}{26} \int \frac{1}{10} \frac{1}{10}$$

Engg I ree.com Steps to design an analog Chebyshew LPF !-1) From the given specifications find the order of the fills N. 2). Round off it to the next higher integer. 3). Wing the following formalas find the values of a and b, which are minor and major and of the ellipse respectively. $a = -\frac{p}{2} \left[\frac{\mu' / w}{2} - \frac{\mu' / w}{2} \right] b = -\frac{p}{p} \left[\frac{\mu' + \mu' / w}{2} \right]$ where $\mu = e^{-1} + \sqrt{e^{-2} + 1} \quad e = \sqrt{10^{0.14} p} - 1$ For normalized Chebyshev fuller sp=Irad/see 4). Calculate the pales of Chebysheo filler which de on an ellipse by using Sk = a cos qk + j 5 sin qk where $q_k = \frac{\pi}{2} + \frac{(2k-i)\pi}{k} = 1, 2, ... N$ 5). Find the denominator polynomial of the Transfer Function using the above podes. 6). The numerator of the Transfer Function depends on the value of N. a) For N-odd, substitute s=0 in D() and find the value. This value is the nelmerator of TF. b). For N-even, substitute s=0 in D(s) and divide the result by 11+02. The value is the numerator of TF.

Design of digital FIR filture -Four techniques! 13 Impulse Invariance (in Bilinear Transformation (iii) Approximation of derivatives. (iv) Matched = transform. Requirements for digital follows to be stable: 1) The jos axes in the s-plane should map into the unit circle in the 2-plane. These atture will be a direct relationship 6/10 the two frequery variables in the two clomeins. (ii) The Lot of the s-plane should map into the inside of the unit circle in the 2-plane. These a stabile analog filter will be convented to a stable digital filler. stype to design digital filters from Analog filters:-D Map the desired digital filter specifications into those for an equivalent analog filter. Derive the analog transfer function for the analog prototype. 3). Transform the transfor function of the analog prototype into an equivalent digital filter transfer Jem ettor. politics show total areas all lines

T1 (107-74) EnggTree.com TIR filter deign by Impulse Invariance Hethy The objective & to design an IIR fills having a unit sample response his, that is the sampled version of the impulse response of the analog filler. ii, h(n) = h(nT) n=0, 1, 2...where T -> sampling interval. From sampling theorem, XLY)= Fo 5 Xa [(f-b) Fs] where f= F/Fs & the normalized frequency. Aliasing occurs if the sampling rate FG des Then twitte the highest frequency contained in 2005 For impulse response, H(f)= Fo 5 Ha [(f-4) Fo] Long $H(\omega) = F_s \stackrel{\neq}{\underset{k=-\infty}{5}} Ha \left[(\omega - 2\pi k) F_s \right] \omega_0$ $H(DT) = \frac{1}{T} \sum_{k=-\infty}^{\infty} Ha\left(D - \frac{2\pi k}{T}\right).$ It is clear that the digital filter with frequency response 1400) has the frequency response characteratics of the corresponding analog feilts if the sampling interval T is schulter sufficiently small, to completely avoid aliasing. If is also clear that impulse invariance method & inappropriate for designing highpass filter due to the spectrum aliasing that results from the sampling process.

EnggTree.com plapping of potute bles z plane and s-plane."-Seotja Zerelu relle e re se e e ST ST IN IN IN IN with rene with parent all the w= AT when o 20 => o 2r 21 0 70 => 1771 0-= 0 => r=1 ii, LHP in spleine & mapped inside the unit the equilited and in the circle in z and the RHP in 5 pleine is mapped outside the unit circle in Z. Also J-2 axes & mapped into the unit circle in 2. However the mapping of the ja axa mes the unit circle & not one-to-one. (Many to one 15-2 - 3TT T s-plane. 2-pilane the for the state of the - 1 -3.4 and

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Since $w \in unique over the range <math>(-\pi, \pi)$ the mapping $w = \alpha \tau$ implies that the interval $-\frac{\pi}{\tau} \leq \pi \leq \frac{\pi}{\tau}$ maps jute the corresponding values of $-\pi \leq \omega \leq \pi$.

The frequency interval TT 5 52 53 th also maps that the interval - Th ZW ET and so on

Let us assume the system function of the analog feller as. in partial fraction. form.

$$Hals) = \frac{s}{k=1} \frac{ch}{s-ph}$$

Also 14

where { ph} are the podes of the analog filter and { ch} are the co-efficients indus pardoas fraction enpansion.

$$halt) = \leq c_k e^{p_k t}$$

 $k=1$ $t \geq 0$

If we sample halt) at tent,

$$h(n) = ha(nT)$$

 $h(n) = \frac{N}{\sum_{k=1}^{N} c_k e^{-p_k T n}}$

i The system function of the resulting digital LTR filter der becomes.

$$H(z) = \sum_{h=0}^{\infty} h(h) z^{-h}$$

$$= \sum_{h=0}^{\infty} \left(\sum_{k=1}^{N} e_k e^{-h(T_h)} z^{-h(T_h)} z^{-h(T_h)$$

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$$= \sum_{k=1}^{N} C_{k} \sum_{z=0}^{n} \left(\frac{e^{P_{k}T}}{1 - e^{P_{k}T}} \right)^{n}$$

$$= \sum_{k=1}^{N} C_{k} \cdot \frac{1}{1 - e^{P_{k}T}} \int \frac{1}{1 - e^{P_{k}T}} \sum_{h \ge 0}^{n} \frac{e^{h}}{1 - e^{h}}$$

$$= \sum_{k=1}^{N} C_{k} \cdot \frac{1}{1 - e^{P_{k}T}} \int \frac{1}{1 - e^{P_{k}T}} \sum_{h \ge 0}^{n} \frac{1}{1 - e^{h}}$$
The agates function of the digital filter e ,

$$H(z) = \sum_{k=1}^{N} \frac{e^{h}}{1 - e^{P_{k}T}} \sum_{z=1}^{n}$$
The digital filter has polle at $z_{k} = \frac{e^{h}}{e^{h}}$
function $H_{a}(z) = \frac{a + e^{h}}{(a + e^{h})^{2} + 9}$
The a digital filter with gutes function $H_{a}(z) = \frac{a + e^{h}}{(a + e^{h})^{2} + 9}$
The a digital filter by mean of the implies incompared method.
Solution $H_{a}(z) = \frac{3 + e^{h}}{(a + e^{h})^{2} + 9} = \frac{8 + 1e^{h}}{3e^{h}} \frac{1}{6 - 2f} + \frac{1}{2}e^{h} \frac{1}{2}f + \frac{$

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$$A = \frac{s+o.i}{s+o.i+j_3} / \frac{1}{s=-o.i+j_3}$$

$$= \frac{-o.i+j_3+o.i}{-o.i+j_3+o.i+j_3}$$

$$= \frac{4x}{-y_4}$$

$$A = \frac{1}{-2}$$

$$B = \frac{-3+o.i}{-y_4+o.i+j_3} / \frac{1}{s=-o.i+j_3}$$

$$= \frac{-p.i-j_3+o.i}{-y_4+o.i+j_3} / \frac{1}{s=-o.i+j_3}$$

$$= \frac{-j_5}{-j_6}$$

$$B = \frac{y_2}{-j_5}$$

$$H = (B) = \frac{1/z}{-j_6} + \frac{1/z}{-j_6} + \frac{1/z}{-j_6}$$

$$Then \quad H(z) = \frac{1/z}{-i-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1}-e^{-i+1$$

For the analog transfer function

$$H(z) = \frac{2}{(z+1)(z+1)}$$

$$H(z) = \frac{2}{z+1} + \frac{2}{z+1}$$

$$H = \frac{2}{z+2} / z = -z$$

$$H = \frac{2}{z+1} / z = -z = z$$

$$H = \frac{2}{z+1} / z = -z = z$$

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$$H = \frac{2}{z+1} / z = -z = z$$

$$H = \frac{2}{z+1} / z = -z$$

EnggTree.com 3) Di analog filter has a transfer finction H(s) = 10 Design a digital filler equivalent to other using impulse invariant method for T=0.2 see, Cliven $H(s) = \frac{10}{s^2 + 7s + 10}$ T = 0.2520. $H(s) = \frac{10}{(s+2)(s+5)}$ $= \frac{A}{S+2} + \frac{B}{S+5}$ $A = \frac{10}{(3+5)} \left(\frac{3+5}{3+5}\right) \left(\frac{3+5}{3-5}\right) = \frac{1}{(3-5)} = \frac{1}{$ $B = \frac{10}{3+2} / 3 = -5 = \frac{10}{-3} = -3.33$ $H(s) = \frac{3 \cdot 33}{3 + 5}$ Form impulse invariance method, For high If HB)= $\sum_{k=1}^{N} \frac{Ck}{s-Pk}$ this $H(z) = \sum_{k=1}^{N} \frac{1-e^{\frac{PkI}{2}}}{k}$ $= \frac{0.666}{0.26782^{4}} = \frac{0.666}{1-0.672^{4}}$ H(2) = 0.20122 1-1.03782 + 0.2472

b) Davign a digities EnggTree.com
the contraint,

$$0.707 \leq \left(H(d^{10})\right) \leq 1$$
, $0 \leq \omega \leq \frac{\pi}{2}$
 $\left(F(d^{10})\right) \leq 0.2$, $\frac{3\pi}{4} \leq \omega \leq \pi$
 $\left(F(d^{10})\right) \leq 0.2$, $\frac{3\pi}{4} \leq \omega \leq \pi$
 $(F(d^{10})) \leq 0.2$, $\frac{3\pi}{4} \leq \omega \leq \pi$
 $(F(d^{10})) \leq 0.2$, $\frac{3\pi}{4} \leq \omega \leq \pi$
 $(F(d^{10})) \leq 0.2$, $\frac{1}{1+d^{2}} = 0.04$
 $e^{2} = 1$, $1+d^{2} = 2.5$
 $e = 1$, $h \in h^{2} = 2.5$
 $e = 1$, $h \in h^{2} = 2.5$
 $e = 1$, $h = h^{2} = 2.4$
 $L d = 0 = 1$, $h = 4.898^{2}$
 $\omega_{p} = \frac{\pi}{2}$, $\omega = 0.5 = \frac{\omega_{1}}{7} = \frac{3\pi}{4}$
 $\frac{3dp-2}{4}$, $\frac{1}{2} = \frac{1}{2} = \frac{1}{2} = \frac{3\pi}{4}$
 $\frac{3dp-2}{1}$, $\frac{1}{2} = \frac{1}{2} =$

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$$\frac{d \sqrt{2}}{d \sqrt{2}}$$
The transfer function H(e) u,
H(s) = $\frac{1}{(s^2 + 0.765s^37 + 0.6s^2 + 1.84773311)}$

S.to-r

aut off frequency.

$$r_{c} = \frac{p}{\sqrt{r_{r}}} = \frac{p}{r_{v}}$$

1

1 ... 10-

· Hab) = Had / 3= 3

: Hald =
$$\frac{(1.57)^4}{(8^2 + 1.2028 + 2.465) (8^2 + 2.9628)}$$

$$H_{2}(y) = \frac{A}{S+1.45+j0.6} + \frac{A^{**}}{S+1.45-j0.6}$$

$$+ \frac{B}{S+0.6+j1.45} + \frac{B^{**}}{S+0.6-j1.45}$$

$$A = 0.7252+j1.754 \qquad B = -0.7252-0.5j$$

$$H_{2}(y) = \frac{0.7252+j1.754}{S-(-1.45-j0.6)} + \frac{0.7252-j1.756}{S-(-1.45+j0.6)}$$

$$+ \frac{-0.7253-0.3j}{S-(-0.6-j1.46)} + \frac{-0.7253+0.3j}{S-(-0.6+j1.46)}$$

$$\text{Engline aca} mathed,
 H(z) = \frac{s}{k_{-1}} \frac{ck}{1-z^{Pk}T_{-1}^{-1}} \quad \text{Telsee.}
 H(z) = \frac{0.7253+j.1.75k}{1-z^{-1/45}z^{-1}j^{-1/55}z^{-1}} + \frac{0.7153-j.1.75k}{1-z^{-1/45}z^{-1}j^{-1/5}z^{-1}} + \frac{-0.7153+0.3j}{1-z^{-0.5}z^{-1}} + \frac{-0.7153+0.3j}{1-z^{-0.5}z^{-1}} + \frac{-0.7153+0.3j}{1-z^{-0.5}z^{-1}} + \frac{-0.7153+0.3j}{1-z^{-1.5}z^{-1}} + \frac{-0.7153+0.3j}{1-z^{-1.5}} + \frac{-0.7153+0.3j}{1-$$

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5) Design a Buttowoorth futtor using implies
nevenance futtor method for the specification

$$0.6 \le | +1(\omega)| \le 1$$
 $0.50 \le 0.23$
 $| +1(\omega)| \le 0.21$ $0.35 \le 0 \le 11$
 $| +1(\omega)| \le 0.21$ $0.35 \le 0 \le 11$
 $| +1(\omega)| \le 0.21$ $0.35 \le 0 \le 11$
 $| +1 = 0.2 \Rightarrow \lambda = 10.89$
 $1 = 0.25 \Rightarrow \lambda = 10.89$
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$$St_{0} = A_{1}^{1} \quad Find \quad add - off \quad Jregundy \quad \underline{Pc}$$

$$= \underline{Pc} = \frac{-2p}{e^{1/N}} = \frac{0 \cdot 257}{(0 \cdot 75)^{1/5}}$$

$$= \underline{Pc} = 0 \cdot 662$$

$$Treation \quad dum (don ::)$$

$$H(3) = \frac{1}{((s+1)(s^{2}+0.6)(s_{0.3}s+1)((s^{2}+1.6)(s_{0.3}s+1))}$$

$$f':=\underline{Pc} = 1 \quad f':=Pc = 1 \quad f'$$

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$$H_{a}(b) = \frac{h}{s+b.6b+} + \frac{b}{st_{0} \cdot s_{0} \cdot s_{1} - j_{0} \cdot b_{1} g} + \frac{b}{st_{0} \cdot s_{0} + j_{0} \cdot b_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1} g} + \frac{c}{s+b.5 \cdot s_{1} + j_{0} \cdot s_{1$$

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$$H(z) = \frac{1 \cdot 2.6}{1 - e^{-0.562} T_{-1}^{-1}} + \frac{-0.09 + \frac{1}{1 - e^{-0.573.6} - \frac{1}{1 - e^{-0.740.7}}}{1 - e^{-0.753.6} - \frac{1}{1 - e^{-0.740.7}}} + \frac{-0.753.6 - \frac{1}{1 - e^{-0.740.7}}}{1 - e^{-0.753.6} - \frac{1}{1 - e^{-0.740.7}}} + \frac{-0.753.6 - \frac{1}{1 - e^{-0.740.7}}}{1 - e^{-0.753.6} - \frac{1}{1 - e^{-0.740.7}}} + \frac{-0.753.6 - \frac{1}{1 - e^{-0.740.7}}}{1 - e^{-0.753.6} - \frac{1}{1 - e^{-0.740.7}}} + \frac{-0.753.6 - \frac{1}{1 - e^{-0.740.7}}}{1 - e^{-0.753.6} - \frac{1}{1 - e^{-0.753.6}}} + \frac{-0.753.6 - \frac{1}{1 - e^{-0.753.6}}}{1 - e^{-0.753.6} - \frac{1}{1 - e^{-0.753.6}}} + \frac{-0.753.6 - \frac{1}{1 - e^{-0.753.6}}}{1 - e^{-0.753.6} - \frac{1}{1 - e^{-0.753.6}}} + \frac{-0.753.6 - \frac{1}{1 - e^{-0.753.6}}}{1 - (0.757.6)} = \frac{1}{1 - e^{-0.753.6}} + \frac{-0.753.6 - \frac{1}{1 - (0.757.6)}}{1 - (0.757.6)} = \frac{1}{1 - e^{-0.753.6}} + \frac{-0.753.6 - \frac{1}{1 - (0.757.6)}}{1 - (0.757.6)} = \frac{1}{1 - e^{-0.753.6}} + \frac{1}{1 - e^{-0.753.6}} = \frac{1}{1 - e^{-0.753.6}} + \frac{1}{1 - (0.757.6)} = \frac{1}{1 - e^{-0.753.6}} + \frac{1}{1 - (0.757.6)} = \frac{1}{1 - e^{-0.753.6}} + \frac{1}{1 - (0.757.6)} = \frac{1}{1 - e^{-0.753.6}} + \frac{1}{1 - e^{-0.753.6}} = \frac{1}{1 - e^{-0.753.6}} =$$

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b Design a chekywa film fir the following
specification wing impute invariance method.

$$D: \xi \leq |H(a|^{0})| \leq 1$$
 $D \leq U \leq 0.2\pi$
 $|H(a|^{0})| = 0 + 0.6\pi \leq U \leq \pi$
 $|H(a|^{0})| = 0 + 0.6\pi \leq U \leq \pi$
 $f_{1+\xi^{2}}$
 $f_{1+\xi^{2}}$
 $f_{1+\xi^{2}}$
 $f_{1+\xi^{2}}$
 $f_{1+\xi^{2}}$
 $h_{2} = 0.2\pi$ $D \leq 0.6\pi$
 $M_{2} = 0.6\pi$
 $M_{2} = 0.6\pi$
 $M_{2} = \frac{(mah^{2}(\frac{1}{2}m))}{(mah^{2}(\frac{1}{2}m))} \leq \frac{(mah^{2}(\frac{1}{2}m))}{(mah^{2}(\frac{1}{2}m))} \leq \frac{(mah^{2}(\frac{1}{2}m))}{(mah^{2}(\frac{1}{2}m))} \geq 1.4\pi$
 $M \geq 1.4\pi$

$$\begin{aligned} & \operatorname{EnggTree.com}_{\alpha = -2p} \left[\frac{p + \frac{1}{2}}{2} \right]_{\alpha = -\frac{1}{2}} \\ &= 0 \cdot 3b \cdot 27 \\ b = -2p \left[\frac{p + \frac{1}{2}}{2} \right]_{\alpha = 0} \\ &= 0 \cdot 7257 \\ \underline{stp > h} \quad Denominator \quad polynomial \quad D(k) \\ &= b \\ \overline{stp > h} \quad Denominator \quad polynomial \quad D(k) \\ &= b \\ &= \frac{\pi}{2} + \frac{(2k-1)\pi}{2N} \quad k = 1, 2, \frac{\pi}{2} \\ &= 1 \\ &= \frac{\pi}{2} + \frac{(2k-1)\pi}{2N} \quad k = 1, 2, \frac{\pi}{2} \\ &= 1 \\ &= \frac{\pi}{2} + \frac{(2k-1)\pi}{2N} \quad k = 1, 2, \frac{\pi}{2} \\ &= 1 \\ &= 0 \cdot 3b + 7 \quad cns (ns^{-1}) + J(0 \cdot 71557) \\ &= 1 \\ &= 0 \cdot 2s \\ &= 0 \cdot$$

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$$\frac{3\frac{1}{2}}{Fyris} = \frac{1}{Fyris} \frac{1}{2} \frac{h}{h(s)} = \frac{k}{k = 1} \frac{c_{s}}{s - h} + \frac{1}{h(s)} \frac{h}{h(s)} = \frac{k}{k = 1} \frac{c_{s}}{s - h} + \frac{1}{h(s)} \frac{h}{h(s)} = \frac{k}{k = 1} \frac{c_{s}}{1 - \frac{h}{h(s)} - \frac{1}{s}}$$

$$= \frac{h}{s + 0 \cdot 25h + j \cdot 513} + \frac{h}{s + 0 \cdot 25h + j \cdot 513}$$

$$A = \frac{0 \cdot 26h}{s + 0 \cdot 25h + j \cdot 513} / \frac{1}{s = -0 \cdot 25h + j \cdot 513}$$

$$A = \frac{0 \cdot 26h}{s + 0 \cdot 25h + j \cdot 513} / \frac{1}{s = -0 \cdot 25h + j \cdot 513}$$

$$B = \frac{0 \cdot 26h}{s + 0 \cdot 25h + j \cdot 513} / \frac{1}{s = -0 \cdot 25h + j \cdot 513}$$

$$E = \frac{0 \cdot 25h}{s + 0 \cdot 25h + j \cdot 513} / \frac{1}{s = -0 \cdot 25h + j \cdot 513}$$

$$B = \frac{0 \cdot 25h}{s + 0 \cdot 25h + j \cdot 513} / \frac{1}{s = -0 \cdot 25h + j \cdot 513}$$

$$E = -0 \cdot 257j$$

$$\frac{1}{s + 0 \cdot 25h + j \cdot 513} - \frac{0 \cdot 257j}{s + 0 \cdot 25h + j \cdot 513}$$

$$\frac{5hp + h}{h(s)} = \frac{0 \cdot 257j}{s + 0 \cdot 25h + j \cdot 513} - \frac{0 \cdot 257j}{s + 0 \cdot 25h + j \cdot 513}$$

$$\frac{5hp + h}{h(s)} = \frac{h}{s + h} + \frac{h(s)}{h(s)} = \frac{h}{k + 1} - \frac{h}{h - e} - \frac{h}{h - e} - \frac{h}{h - e} - \frac{1}{h - e} - \frac{h}{h - e} - \frac{1}{h - e} - \frac{h}{h - e} - \frac{1}{h - e} - \frac$$

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$$f(z) = \frac{(-2\pi)^2}{(-2-6\pi)^2} \frac{(-2\pi)^2}{(-2\pi)^2} \frac{(-2\pi)^2}{(-2-6\pi)^2} \frac{(-2\pi)^2}{(-2-6\pi)^2} \frac{$$

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$$\frac{d\eta(H)}{dt} + a \ \eta(H) = bx(H) \longrightarrow 0$$
The type that the derivative and approximate the integral by the tropezoidal dormale.

$$\therefore \ \eta(H) = \int_{0}^{T} \int_{0}^{T} (t) \ dt + topezoidal dormale.$$

$$\therefore \ \eta(H) = \int_{0}^{T} \int_{0}^{T} (t) \ dt + topezoidal dormale.$$

$$dt = \int_{0}^{T} \int_{0}^{T} (t) \ dt + topezoidal dormale.$$

$$dt = nT \ and \ t_{0} = nT - T \ dt = b \ dt + a \ dt = b \ dt \ dt = b \$$



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If rei then 0 >0

in the LHP of splane maps inside the with circle. and RHP of splane maps outside unit circle.



The entire range in s & Mapping between two mapped only once into this frequencies.

sontineer. But the mapping to highly

We observe a frequency compression was frequency warping as it is usually called, due to the nonlinearity of ethe arctangent function.

The bilinear transformation maps the point $s = \infty$ into the point z = -1.

$$\begin{array}{l} \begin{array}{l} \begin{array}{l} \text{Apply Billinear transform } t_{1} \\ \begin{array}{l} \text{Engine result} \\ \text{P}(5) = \frac{1}{((z+1)(3+3)} \\ \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \text{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \mbox{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \mbox{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \mbox{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \mbox{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \mbox{Weight } \\ \mbox{Weight } \\ \end{array} \end{array} \\ \begin{array}{l} \text{Weight } \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \end{array} \\ \begin{array}{l} \text{Weight } \\ \mbox{Weight } \\ \mbox{Weight$$

O. Design a digital Butterworth feller saterfying the following specification 0.8 ≤ [H[20)] ≤1 0 ≤ [6] 3 0.20 (+(e))) SOL 0.6TT S 10/STT Assume TESP SEC. Apply Bilinear transformation Sole Given 1 = 0.8 = 5 = 0.75 $\frac{1}{1+\lambda^2} = 0.2 \implies \lambda = 4.898$ $\omega_p = 0.2\pi \quad \omega_s = 0.6\pi$ Stip-1: From Bilinear transformation, $-\frac{1}{2p} = \frac{2}{T} \tan \frac{w_p}{2} = \frac{2}{D} \tan \frac{0.2\pi}{2} = \frac{1}{0} - \frac{1}{2} - \frac{1}{2}$ as = 2 fan ws = B / A / 0 27-52 2.752 1.013 8 tip- 2 Analog filter Transfer Function Haw step 2.1 Order N $N \ge \frac{\log\left(\frac{x}{c}\right)}{\log\left(\frac{-\Omega_{s}}{\exp}\right)}$ $\geq \frac{\log\left(\frac{4\cdot 191}{0.15}\right)}{\log\left(\frac{2.7152}{0.64.95}\right)} \geq \frac{0\cdot 5\cdot 149}{0.62.65}$ N Z 1:29 Step 2.2 Round-off N to the next higher inleger : [N=2]

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$$stp = 1 \cdot 3 \qquad \text{Analog TF H(s) for } c_c = 1 \text{ real bace}$$

$$p(S) = s^{2} + f = s + 1$$

$$H(s) = \frac{1}{D(S)} = \frac{1}{s^{2} + f = s + 1}$$

$$stp = 2 \cdot 4 : \quad Cud = off \qquad frequency \quad P = C$$

$$-Cc = -\frac{2 \cdot h}{(c)^{1/n}} = -\frac{0!64}{(c \cdot 7s)^{1/2}}$$

$$= 0 \cdot 70 \cdot 5 \circ 3$$

$$stp = 1 \cdot 7 \qquad \text{Analog TiF. } H_{2}(S)$$

$$H_{0}(S) = -\frac{1}{(c)^{2} + (c \cdot 2)^{2}}$$

$$= -\frac{1}{(c)^{2} + (c \cdot 2)^{2}}$$

$$= -\frac{1}{(c)^{2} + (c \cdot 2)^{2}}$$

$$H_{0}(S) = -\frac{1}{(c \cdot 2)^{2}}$$

$$= -\frac{(c)(7 \cdot 6 \cdot 3)}{(c \cdot 7s)^{2}} + \frac{1}{(c \cdot 2)^{2}}$$

$$H_{0}(S) = -\frac{0 \cdot 5 \cdot 4 \cdot 2}{s^{2} + U \cdot 6 \cdot 16t \cdot 3} + 0 \cdot 5 \cdot 5 \cdot 17$$

$$= -\frac{(b \cdot 5 \cdot 4 \cdot 2)}{(c \cdot 7s)^{2}}$$

$$H_{1}(S) = -\frac{0 \cdot 5 \cdot 4 \cdot 2}{s^{2} + U \cdot 6 \cdot 16t \cdot 3} + 0 \cdot 5 \cdot 5 \cdot 17$$

$$= -\frac{(c \cdot 5 \cdot 4 \cdot 2)}{(c \cdot 7s)^{2}} + 1 \cdot 0 \cdot 10t \cdot 5 \left(\frac{1}{7} + \frac{1 - 2^{2}}{(1 + 2^{2})}\right) + 0 \cdot 5 \cdot 5 \cdot 25$$

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EnggTree.com $0 \cdot 5629 (1+z^{-1})^{2}$ $H (1+z^{-1})^{1} + 2(1+608)(1+z^{-1})(1+z^{-1})+0 = 5449 (1+z^{-1})^{2}$ $0 \cdot 5629 (1+z^{-1})^{2}$ $H - 8z^{-1} + 4z^{-2} + 1\cdot 2z16 - 1\cdot 2z16z^{-1} + 0 \cdot 5629 + 1\cdot 1258 z^{-1} + 2x7825 = 0.5629 = 0.5629 (1+z^{-1})^{2}$ $= \frac{0 \cdot 5629 (1+z^{-1})^{2}}{5 \cdot 8745 - 5629 (1+z^{-1})^{2}} + 2 \cdot 7784$ $H (z) = \frac{0 \cdot 0958 (1+z^{-1})^{2}}{1 - 1 \cdot 07435z^{-1}} + 0 \cdot 47295z^{-2}$

3. Design a Chebysher filter for the following specification using bilinear finingermation $0.8 \le [H[e]^{(0)}] \le 1$ $0 \le 10 \le 0.25$ $[H[e]^{(0)}] \le 0.2$ $0.657 \le 10 \le 17$ (lines) $\frac{1}{1+e^2} = 0.8 \Rightarrow 6 = 0.75$ $0 \ge 0.25$ $\frac{1}{1+e^2} = 0.2 \Rightarrow \lambda = 1.898$ $\int 1 + k^4$ Let $T = 1 \le 2$ $\int 1 + k^4$ $D = 2 \le 100$ $\frac{0.25}{2}$ D = 0.248

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$$\begin{aligned}
\left(\nabla_{2} = \frac{\pi}{2} + \frac{3\pi}{4} = 2^{2}\pi^{*} \\
S_{k1} = 0.3752 \cos(135^{\circ}) + J(0.75) \cos(135^{\circ}) \\
S_{1} = -0.2653 + J0.53 \\
B_{2} = 0.8752 \cos(2.25^{\circ}) + J(0.75) \cos(2.25^{\circ}) \\
S_{2} = -0.2653 - J0.53 \\
S_{2} = -0.2653 - J0.53 \\
S_{2} = -0.2653 - J0.53 \\
D(3) = (3 + 0.52654 + 0.3516 \\
S_{2} = \frac{0.3516}{(1 + e^{2})} = \frac{0.3516}{(1 + 0.75^{2})} \\
N_{T} = \frac{D(3)}{3 = 0} = \frac{0.3516}{(1 + 0.75^{2})} \\
N_{T} = 0.27 \\
N_{T} = 0.27 \\
N_{T} = 0.27 \\
H_{a}(3) = \frac{0.2}{3^{2} + 0.53068 + 0.3516} \\
S_{2} = \frac{0.35}{(1 + 2^{2})^{2}} \\
= \frac{0.2}{\pi} \frac{1 - 2^{-1}}{(1 + 2^{2})} \\
= \frac{0.2}{\left[2 \frac{1 - 2^{-1}}{(1 + 2^{2})}\right]^{2} + 0.5306\left[2 \frac{1 - 2^{-1}}{(1 + 2^{2})}\right] + 0.35^{14}
\end{aligned}$$

$$\begin{aligned} & \text{ErggTree.com} \\ & \text{H}(s) = \underbrace{\begin{array}{c} 0 \cdot 18 \left((1+z^{-1})^{2} \\ \hline 5 \cdot 11 28 - 7 \cdot 29 8 z^{-1} + 3 \cdot 29 z^{-2} \end{array}}_{5 \cdot 11 28 - 7 \cdot 29 8 z^{-1} + 3 \cdot 29 z^{-2}} \\ \hline \\ & \text{H}(z) = \underbrace{\begin{array}{c} 0 \cdot 052 \left((1+z^{-1})^{2} \\ \hline 1 - 1 \cdot 348 z^{-1} + 0 \cdot 608 z^{-1} \end{array}}_{1 - 1 \cdot 348 z^{-1} + 0 \cdot 608 z^{-1}} \end{aligned}} \\ \\ & \text{Dre} \quad \text{Of lite Design by Popposimation of Derivative}}_{an analog fulter into a digital futter is to a analog fulter into a digital futter is to a opproximate the differential equivation of the differential equivation and the differential equivation is the differential equivation is the differential equivation of the differential equivation is the difference is$$

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The analog differentiator with output

$$\frac{dy(t)}{dt}$$
 has the system function $t(a) = d$
 $\frac{dy(t)}{dt}$ has the system that produce the
 $output [\underline{y}(n) - \underline{y}(n-1)]$ has the system function
 $\mu(a) = (\underline{1-z^{-1}})$.
 $T = second derivative $\frac{d^2y(b)}{dt^2}$, is replaced
by the second difference, which is derived
as follows:
 $\frac{d^2y(t)}{dt^2}/t = nT = \frac{d}{dt} \cdot \left[\frac{d}{dt}\frac{dt}{dt}\right]_{t=nT}$
 $= \frac{y(n) - 2y(n-2)}{T} - \left[\frac{y(n) - 1}{T}\right] - \left[\frac{y(n) - 1}{T}\right]$
 $T = \frac{y(n) - 2y(n-2)}{T} - \frac{y(n-2)}{T}$
 $T = \frac{y(n) - 2y(n-2) + y(n-2)}{T^2}$
 $T = \frac{(1-2z^{-1})^n}{T} = 0$$

EnggTree.com . The system function for the digital IIR filter obtained as a result of the approved matters of the derivatives by finite difference & $H(=) = H_{a} (u) / u = \frac{1-u}{u}$ Where Hals) -> Analog falter syster functions. pollin an has Mapping : $Z = \frac{1}{1 - sT}$ $S = \frac{1}{3}P_{1}, \quad Z = \frac{1}{1 - \frac{1}{3}P_{1}T}$ $= \frac{1}{1+2^{2}T^{2}} + \frac{1}{1+2^{2}T^{2}}$ As 52 varies from - as to as, the corresponding local of points in the z-plane is a circle of raching 1/2 and with center at 2=1. Under the 2-plane 8-pilane The mapping is 1-2t takes LHP in the 8-plane into points inside the circle of radius 1/2 and center z=1 in the z-plane.

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The mapping takes points in the LHP of the 3-plane into corresponding points inside the dirde in the 2-plane, and points in the RHP of the s-plane are mapped into points outside the circle.

- "I The mapping has a destrolate property filles The mapping has a destrolate property filles Their a stable ancience of the digital fill filter and to are confirmed to relatively small frequencies and as a consequence, this mapping to restricted to the disign of LPF and BDF houring relatively small resonant frequences
- * It is not possible for enample to transform a signpois analog filler into a corresponding highpass digital filler.

() Convert the analog bandpass feller with system function

Hale): (3+0-1) + 9 Into a digital IJE falter by use of the backwood difference for the derivatives.

(given, Hals) = 1 (0+0.1)2+9 S= 1-2-1

+(=)= --

 $= \frac{1}{(1-z^{-1})^2 + 2(1-z^{-1})^{0.1} + 0.1^2 + 9}$ Downloaded from EnggTree.com

 $\int \frac{1-z^{-1}}{\tau} + 0 \cdot 1 + q \cdot$

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$$= \frac{\tau^{2}}{1 - 2z^{2} + z^{2}} + \vartheta \cdot 2 - 0 \cdot 2z^{2} + \vartheta \cdot 0 + 9}$$

$$= \frac{\tau^{2} \int (\mu \circ \cdot 2\tau + 9 \cdot \circ 1\tau^{2})}{1 - \frac{2(1+0\cdot17)}{1+0\cdot2\tau + 9 \cdot \circ 17}} + \frac{1}{1+0\cdot2\tau + 9 \cdot \circ 17} = z^{2} \int (\mu \circ \cdot 2\tau + 9 \cdot \circ 17)$$
If have the Yorm of resonator provided that τ
is selected small enough $(T \leq 0 \cdot 1)$ in order for the
poly is be near the unit circle. The condition
 $q_{1}^{2} \leq h \circ 2$ is satisfied, so that the poly one
 $q_{1}^{2} \leq h \circ 2$ is satisfied, so that the poly one
 $q_{1}^{2} \leq h \circ 2$ is satisfied, so that the poly one
 $(\sigma - y) = -\sum_{k=1}^{N} a_{k} \int (1-k) + \sum_{k=0}^{H} b_{k} \cdot z(n-k)$
 $H(z) = \frac{y}{1+\sum_{k=1}^{N} a_{k}} - 0$
Different Tit futter realizations are given based
on age 0 and Co.
Types of structure:
 0 Direct form $S \geq 1$
 0 Direct form $S \geq 1$

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Direct sports Stuctures ?





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Parallel form structure :-

C

A parallel form realization of an 11k system can be obtained by performing a partiel fraction expansion of H(=). With old closs of generality, assume that NZH and that the poly are defined.



Subsystims has the form, $t_{k}(z) = \frac{b_{k0} + b_{k1} z^{-2}}{1 + a_{k2} z^{-2}}$











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$$\begin{aligned} & \text{EnggTree.com} \\ & \text{where } \oint U_m \oint \text{ are due parameters} & \text{thet determine} \\ & \text{Hermore of the Aythirs.} \\ & \text{H}(z) = \frac{Y(z)}{x(z)} \\ & = \frac{M}{z(z)} & \text{Un}\left(\frac{Um(z)}{x(z)}\right) \\ & = \frac{M}{m_{2D}} & \text{Un}\left(\frac{Um(z)}{x(z)}\right) \\ & \text{Hermore for an equation of the equation of the$$

For m=3, k=1

$$a_{2}(D = \frac{a_{3}(D - a_{2}(3) a_{3}(2)}{1 - a_{3}^{2}(3)}$$

$$= \frac{\frac{13}{24} - \frac{1}{3} \left(\frac{57}{8}\right)}{1 - \left(\frac{2}{3}\right)^2} = \frac{3}{8}$$

For m=3 k=2

$$k_2 = a_2(2) = \frac{a_3(2) - a_3(3) a_3(1)}{1 - a_3^2(3)}$$

$$=\frac{5}{6}-\frac{1/3}{\left(\frac{13}{2h}\right)}=\frac{1}{1-\left(\frac{1}{3}\right)^2}=\frac{1}{2}$$

For
$$m=2$$
, $k=1$
 $k_1 = q_1(1) = \frac{q_2(1) - q_2(2) q_2(1)}{1 - q_2^2(2)}$

$$= \frac{3/8 - \frac{1}{2} \left(\frac{3}{2}\right)}{1 - \left(\frac{1}{2}\right)^2} = \frac{1}{4}$$

For Ladice structure,

$$k_1 = \frac{1}{4}$$
 $k_2 = \frac{1}{2}$ $k_3 = \frac{1}{3}$

For Ladder structure $\begin{array}{c} M\\ C_{m}=b_{m}-\frac{5}{2} \quad Ci\,ai\,(i-m) \quad m=M,\,M-1\dots 0\\ i=m+1 \quad C_{3}=b_{3}=1\\ C_{2}=b_{2}-C_{3}a_{3}(1)\\ =2-1\left(\frac{13}{24}\right)=1\cdot4583\\ C_{1}=b_{1}-\frac{5}{2} \quad c_{3}a_{1}(1-m)\\ =b_{1}-\int c_{2}a_{2}(1)+c_{3}a_{3}(1) \end{bmatrix}=0\cdot B_{2}F1\\ Downloaded from EnggTree.com\end{array}$

$$c_{0} = b_{0} - \sum_{i=1}^{3} c_{i} a_{i} \left(1 - m\right)$$

$$= b_{0} - \left[c_{1} a_{i} (1) + a_{2} a_{2} (z) + c_{3} a_{3} (z) \right]$$

$$= - 0 \cdot 2^{15} \pi$$

$$The Lattice Ladder stracture b,$$

$$\frac{z^{(n)}}{y^{(n)}} + \frac{y^{(n)}}{y^{(n)}} +$$

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$$Childy hw Fills Design
Order N:
$$[H(j, a)]^{2} = \frac{1}{1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)} \qquad Al = 1, 2, ..., \\
H(j, a)]^{2} = \frac{1}{1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)} \qquad -0$$

$$Taking (al-eg) \quad aboth sides,
so deg [H(j, a)] = lo deg 1 - lo log [1 + e^{2} C_{N} \left(\frac{a}{ap}\right)]$$

$$^{1} 2o deg [H(j, a)] = -10 \log [1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

$$^{2} 2o deg [H(j, a)] = -10 \log [1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

$$^{2} 2o deg [H(j, a)] = -10 \log [1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

$$^{2} 2o deg [H(j, a)] = -10 \log [1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

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$$^{2} 2o deg [H(j, a)] = -10 \log [1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

$$^{2} 2o deg [H(j, a)] = -10 \log [1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

$$^{2} 2o deg [1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

$$^{2} ce^{2} = 10^{0.1 de_{p}-1} \qquad (3)$$

$$^{2} def = 10 \log [1 + e^{2}] / \cdots C_{N} (D = 1)$$

$$^{2} e^{2} = 10^{0.1 de_{p}-1} \qquad (3)$$

$$^{2} le^{2} def (1 + e^{2} C_{N}^{2} \left(\frac{a}{ap}\right)]$$

$$^{2} def def attenuation in step band ba ds.$$

$$^{2} des = +10 \log [1 + e^{2} (cash (N cosh^{-1} \left(\frac{a}{ap}\right))]$$

$$^{2} des = 10 \log [1 + e^{2} (cash (N cosh^{-1} \left(\frac{a}{ap}\right))]$$

$$^{2} des = 10 \log [1 + (e^{0.1 de_{p}-1)] so sh (N cosh^{-1} \left(\frac{a}{ap}\right)]$$

$$^{3} lo^{0.1 de_{p}} -1 = (10^{0.1 de_{p}-1}) so sh (N cosh^{-1} \left(\frac{a}{ap}\right)]$$$$

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$$\frac{\mu_{p}}{\mu_{p}} = \frac{1}{1} = \int \cosh \left[N \cosh V^{+} \left(\frac{m_{p}}{m_{p}} \right) \right]^{2}$$

$$N \cosh V^{-1} \left(\frac{m_{p}}{m_{p}} \right) = \cosh V^{-1} \left[\frac{\mu_{p}}{\mu_{p}} \frac{\mu_{p}}{\nu_{p}} - \frac{1}{1} \right]$$

$$N = \frac{(2\pi)}{(2\pi)^{2} (2\pi)^{2} (2\pi$$

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$$\begin{aligned}
\pm \int_{e} = \cos \left[N \varphi - j N \varphi \right] \\
= \cos N \varphi \cos \left[N \varphi + \sin \left(N \varphi \right) \right] \wedge \left(j N \varphi \right) \\
= \cos \left(N \varphi \right) \cos h \left(N \varphi \right) + j \sin N \varphi \sin h \right] \\
= \cos \left(N \varphi \right) \cos h \left(N \varphi \right) + j \sin N \varphi \sin h \right] \\
Equating read and imighting ports,
$$\cosh \left[N \varphi \right] \cosh \left[N \varphi \right] = 0 \qquad (3)$$
Sin $\left[N \varphi \right] \cosh \left[N \varphi \right] = 0 \qquad (4)$
Sin $\left[N \varphi \right] \cosh \left[N \varphi \right] = 0 \qquad (4)$
Sin $\left[N \varphi \right] \sinh \left[N \varphi \right] = \frac{1}{2\pi} = -\varphi$
Sin $(N \varphi) \sinh \left[N \varphi \right] = \frac{1}{2\pi} = -\varphi$
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Sin $($$$

$$g_{Nh} k = g_{Nh} \left[\frac{1}{\mu} \frac{g_{Nh} h^{2}(e^{-1})}{g_{Nh} h^{2}(e^{-1})} \right]$$

$$= \frac{e^{\frac{1}{\mu}g_{Nh} h^{2}(e^{-1})}}{e^{-\frac{1}{\mu}(e^{-1})}} \frac{1}{\mu}$$

$$= \left[\frac{e^{\frac{1}{\mu}g_{Nh} h^{2}(e^{-1})}}{e^{-\frac{1}{\mu}(e^{-1})}} \frac{1}{\mu} - \frac{1}{\mu} \frac{1}{\mu} \frac{1}{\mu} - \frac{1}{\mu} \frac{1}{$$

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EnggTree.com FIR FILTER DESIGN Structures of FIR - Linear phase FIR filler -Fourier servis - Filter design using windowing techniques [Rectangular windows, Hamming window, Hanning window] Frequency sampling techniques - Fruite word length effects in digital fillers: Errors, Linet Cycle, Noae power spectrum. antitale the cores of the 4 INTRODUCTION ? TI (660-664) FIR filters are non-recersive type. The present output depends on present and precious values of input. $y(n) = b_0 x(n) + b_1 x(n-1) + \cdots + b_{M-1} x(n-M+1)$ $g(n) = \sum_{k=0}^{M-1} b_k r(n-k)$... $y(z) = \sum_{k=0}^{M-1} b_k z^{-k} x(z)$ $H(2) = \frac{Y(2)}{x(2)} = \frac{M-1}{k=0} b_{k} = \frac{-k}{k}$ Linear phase FIR fillers: An FIR filter of length M & given by extre @ and @. where I but set of filter co-efficients. Alternatively, the output sequence as its the consolution of the cenif sample response h(n) of the system with the input signal.

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H-1

$$f(n) = \frac{H}{k_{0}} h(k) z(n-k)$$

Equal Q and Q are identical.
Equal Q and Q are identical.
The filter can also be characterized by it system for
 $H(z) = \frac{1}{k_{0}} h(k) z^{-k}$
 $h(z) = \frac{1}{k_{0}} h(z) \int z(z) \int z($

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must occur in EnggTree compairs. i, If =, & a root cors a zero of H(2), then 1/2, & also a root. If due unit sample response him of the filter & real, complex valued roots must occurs In complex - conjugate pairs. Hence, If ZI & a complex - valued root, Zit & also a root. As a consequence of (F), H(=) also has a zero at 1 . · Vzi 0 23 to wit arde. 121 23 0 + et. 1/231 Fig. symmetry of zero locations for a linear phase FIR filler. The frequency response characteristics of linear-phase FIR fillers are obtained by evaluating (B) on the unit arde. The substitution yields the expression When h(n) = h [M-1-n], H(w) can be expressed as for H(w). $H(\omega) = H_r(\omega) = -j \omega (M-D/2)$ 8 where Hr (w) is a real function of us and can be $H_{T}(w) = h\left(\frac{M-1}{2}\right) + 2\frac{5}{2} h(w) \cos w\left(\frac{M-1}{2} - 5\right), \mod 1$ expressed as $H_r(w) = 2 \frac{5}{22}$ $h(n) \cos \left(\frac{M-1}{2} - n\right)$, Mever

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The phase characteristic of the filter for book
H-odd and Heven is,

$$\int_{-\infty}^{+\infty} \left(\frac{H-1}{2}\right) + \pi \quad \text{if } H_{T}(\omega) > 0$$
(i) $\omega = \int_{-\infty}^{+\infty} \left(\frac{H-1}{2}\right) + \pi \quad \text{if } H_{T}(\omega) < 0$
(i) $\omega = \int_{-\infty}^{+\infty} \left(\frac{H-1}{2}\right) + \pi \quad \text{if } H_{T}(\omega) < 0$
(i) $\omega = \int_{-\infty}^{+\infty} \left(\frac{H-1}{2}\right) + \pi \quad \text{if } H_{T}(\omega) < 0$
(i) $\omega = \int_{-\infty}^{+\infty} \left(\frac{H-1}{2}\right) = 0$
How even, if the easter point of the anti-symmetric.
 $h(n) \in n = (H - n)/2$, consequently,
 $h\left(\frac{H-1}{2}\right) = 0$
How even, if $M \in \text{even, each then fn h(n) has a matching term of opposite sign.
The forgunary response of an File fills with an actage metric with sample response con be expressed as,
 $H(\omega) = H_{T}(\omega) = \int_{-\infty}^{+\infty} h(n) \sin \omega \left(\frac{H-1}{2} - n\right), m - odd$
 $\frac{H-1}{2 - n \geq n} = h(n) \sin \omega \left(\frac{H-1}{2} - n\right), M - \omega m$
The phase charactivistic of the filter for both
 $M - add$ and M even is
 $M - add$ is $(M - \frac{1}{2}), \quad \text{if } H_{T}(\omega) > 0$
For a symmetric $h(n),$
the number of fills is -off their specify the formous,
 $M - \frac{1}{M} = 0$ or $M + 1$ and $M = 0$.
If the unit sample response to antisymmetric,
 $h\left(\frac{M-1}{M}\right\right) = 0$ M is only metric,
 $H\left(\frac{M-1}{M}\right) = 0$ M is only metric.$

DESIGN OF LINEAR PHASE FIR FILTERS USING WINDOWS The method of design starts with the desired frequency response specification Hd(w) and the corresponding unit sample response hallos & determined. Hollow) = 5 holon _____ where hd (s) = 1 | Hd (co) e' word ou = 3 Thus given Hodow), we des determine the unit sample response hollow by evaluating the integral in @ The unit sample response halos obtained from () is infinite in duration and must be truckled at some point, say at n=M-1 to yield an FIR filter of length M. Tradition of hollow to a length M-1 & equivalent to multiplying holling a "rectangular window" defined as, $w(m) = \begin{cases} 1, & n = 0, 1, \dots, M-1 \\ 0, & other wise \end{cases}$ - (3) \therefore h(n)=hd(n) w(n) = { hdin n=0,1,...M-1 otherwise The multiplication of hd(n) with window function a equivalent to consolution of Adlos) with Wlas) where $W(w) = \frac{M-1}{5} w(n)e^{-jwn}$ Thus the convolution of Adw with Wolco) yields the fraquency response of the (trandleted) FIR filler.

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$$w$$
, $H(w) = \frac{1}{2\pi} \int_{-\infty}^{\infty} H_{d}(w) W(w - w) dw$
The Fourier Transform of readongalar wholes
 a , $W(w) = \sum_{n=0}^{\infty} e^{-jwH}$
 $= \frac{1-e^{-jwH}}{1-e^{-jw}}$
 $= e^{-jw(H-1)/e} \frac{\sin(wH/e)}{\sin(w/e)}$
The window function has a magnitude response.
 $|W(w)| = \frac{(3m(wH/e))}{|sin(w/e)|}$, $\pi \le w \le \pi$
and $O(w) = \int_{-\infty}^{\infty} \frac{(H-1)}{|sin(w/e)|} + \pi$ when $\sin(w/He) \ge 0$
The magnitude response of the window function
 w ,
 $\int_{-\infty}^{\infty} \frac{(H-1)}{|w|} + \pi$ when $\sin(w/He) \ge 0$
The magnitude response of the window function
 w ,
 $\int_{-\infty}^{\infty} \frac{(H-1)}{|w|} + \pi$, $h \ge H$ increases
the width will be come neartower. However the
sidebalas of $|W(w)|$ one relating by and wolffield
by an increase in H.

* The characteroities of rectangular window play a significant rode in determining the resulting frequency rosponse of the FIR filter obtained by trancating halfs to length M. A The convolution of Holes) with W(w) has the effect of smoothing Hdlw). As MEs increased, W(w) becomes norrower, and the smoothing provided by w (w) & reduced. * On the otherhoard, the large side do bes of W(2) result in some undisirable ringing effects in the Fir filter frequency response HCWD, and also 30 relatively larger side lobes in H (w). * These undesirable effects are best allevided by the use of windows that do not contain abrupt des continuitées in their line - domain characterenties and have correspondingly low side tobes 90 their frequency-domain characteredics. Window Functions for FIR filler Design Tince domain sequence Xance of Window him, osneH-1 2 n-<u>H-1</u> O Bartlett (Triangular 1 - - M-0.42 - 0.5 cos 270 +0.08 cos 400 M-1 2). Blackman 0.54-0.46 cos 210 M-1 3) Hamming $\frac{1}{2}\left[1-\cos\frac{2\pi h}{M-1}\right]$ 4) Hanning

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Clibbs phenomenon:-

The multiplication of hdm with a rectargular window & Fdentical to truncating the Fourier serves representation of the desired filler characters, Hollow). The francation of the Fourier serve & known to introduce ripple in the frequency response characters to H (w) due to the nonuniform convergence of the Fourier serves at a descontinuity. The oscillatory behavior near the band edge of the Fourier serves of filler & called the Gibbs phenomenon.

Inportant Frequency - Domain Characteratics of Window Functions		
Type of Window	Approximate transition width of main lobe	Peak side lobe (dB)
0 Dectangular	<u>477</u>	-13
D artilett	BX/M	- 25
2) Burning 3) Hanning	8π/M 8π/M 12π/M	-31 -41
4) Hamming		-57
5) Bleek		down H The

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$$sbg^{-1}: \quad Frequency \quad esponse
H(e)^{(0)} = \int_{\pi=0}^{4} h(x) e^{-\int_{\pi}^{(0)} x^{-1}} + h(y) e^{-\int_{\pi}^{(0)} x^{-1}} + \frac{h(y) e^{-\int_{\pi}^{(0)}$$

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27

> w->

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$$Hd \left(e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac{1}}}e^{\frac{1}{1}}e^{\frac{1}{1}}e^{\frac$$

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Window Function,

$$\omega(w) = \int_{0}^{1} \int_{0}^{1} \frac{\omega(w)}{\omega h^{\omega}}$$

 $file = \int_{0}^{1} \int_{0}^{1} \frac{\omega(w)}{\omega h^{\omega}}$
 $h(w) = hd(m) \omega(m)$
 $h(w) = hd(m) \omega(m)$
 $h(w) = hd(m) \omega(m)$
 $h(w) = 0.068w = h(4)$
 $h(1) = \frac{\omega(w)}{\omega m^{\omega}} = h(4)$
 $h(3) = 0.4$
 $file = \frac{1}{2}$
 fi
$$= \frac{1}{2\pi} \int_{-\infty}^{\infty} \frac{e^{-j\pi(n-2)}}{j(n-2)} \frac{\pi}{-\pi} + \int_{-\infty}^{\infty} \frac{e^{-j\pi(n-2)}}{j(n-2)} \frac{\pi}{-\pi} \frac{\pi}{-\pi} \int_{-\infty}^{\infty} \frac{e^{-j\pi(n-2)}}{j(n-2)} \frac{\pi}{-\pi} \frac{\pi}{-\pi} \int_{-\infty}^{\infty} \frac{e^{-j\pi(n-2)}}{j(n-2)} \frac{\pi}{-\pi} \int_{-\infty}^{\infty} \frac{e^{-j\pi(n-2)}}{j(n-2)} \frac{\pi}{-\pi} \int_{-\infty}^{\infty} \frac{e^{-j\pi(n-2)}}{j(n-2)} \frac{\pi}{-\pi} \int_{-\infty}^{\infty} \frac{\pi}{-\pi} \int_{-\infty}^{\infty} \frac{e^{-j\pi(n-2)}}{j(n-2)} \frac{\pi}{-\pi} \int_{-\infty}^{\infty} \frac{\pi}{-\pi} \int_{-\infty$$

EnggTree.com Step-5 FIR filler co-efficients hen = hden when m $h(0) = hd(0) \ \omega_{H_{A}}(0) = \Theta \frac{1}{2\pi} |0| = 0 = h(4)$ $h(i) = hd(i) \cup An(i) = \frac{1}{1 - \pi} (0.\pi) = 0.1125 = h(3)$ $h(2) = hd(2) W H h(2) = \frac{3}{4} (1) = \frac{3}{4} = 0.75$

4). Design a digital FIR band - pass filter with lower cut-off frequency 2000Hz and upper cut off frequency 3200Hz using Hamming window of length N=7. Sampling rate & 10000 Hz. (NHD-12)

Solution :-

Cliven, Filter Type: Band Pass Filter

Lower autoff fraguency : 2000 HZ Upper cutoff frequency: 3200 Hz Sampling rate : 10000 Hz Length M: 7 : Hamming Window window

step-1 hdln): $k_{\partial} + d(e^{j\omega}) = \begin{cases} e^{-j\alpha \omega} & \omega_{c_{1}} \leq |\omega| \leq \omega_{c_{1}} \\ 0 & \omega_{c_{2}} \leq |\omega| \leq \pi \end{cases}$ ∞wc, ≤/w/≤wc2 OLIDIENG $\alpha = \frac{M-1}{2} = \frac{6}{2} = 3$ $w_{c_1} = \frac{2 \times \pi + c_1}{F} = \frac{2 \times \pi \times + 900}{600} = \frac{2\pi}{5} = 0.4\pi$ $w_{c_{1}} = \frac{2\pi + e_{2}}{F} = \frac{\beta_{x} \pi \times 3200}{10,000} = 0.64 \pi$

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DESIGN OF LINIEAR PHASE FIR FILTERS BY INE FREQUENCY - SAMPLING METHOD TI (671-673)

In the frequency -sampling method for FIR filter durger, we specify the desired frequency response Hd(w) at a set of equally spaced

frequencies, namely $w_k = \frac{2\pi}{M} (k + \alpha), \quad k = 0, 1, \dots \frac{M-1}{2}, Modd$

 $k=0,1,\cdots,\frac{M}{2}-1$, Meven

and solve for the unit sample response her of the FIR filter from these equally spaced frequency specifications. To reduced side doses, it is desirable to optimize the frequency specifications in the transition band of the filter.

A basic symmetry property of the sampled frequency response function to simplify the computation is given her set us begin with the desired frequency response of the FIR filler, H-1 herr e-Joon H(w) = 5 herr _____

If we specify the frequency response of the filter at the frequencies given by O, then from O $H(k+\alpha) \equiv H\left(\frac{2\pi}{H}(k+\alpha)\right)$

 $H\left[k+\alpha\right) = \frac{M-1}{2} h(n) e^{-j2\pi(k+\alpha)n/m}$ k= 0,1. . - M-1

EnggTree.com It a simple to invert @ and express h(h) interms of H(k+x). If we multiply both sides of (3 by the exponential, exp (j2th hm/m), m=0, 1, ... M-1 and sum over k=0, 1,... M-1, the right hand side of 3 reduces to them M h[m) exp[-j24xm/H). $L(h) = \frac{M-1}{M} = \frac{M-1}{K+\alpha} = \frac{2\pi (k+\alpha)n/m}{M-k=0}$, n=0,1,...M-1 -4 From above set a, we compute the values of the unit sample response h(n) from the specification of the frequency samples H(let a), k=0,1,...,H-1. When \$20, equils reduces to the DFT of the sequence Shirs? and equile reduces to IDFT. Since him is real, we can easily show that the trequency samples } H(k+a) } satisfy the symmetry condition, $H(k+\alpha) = H^{*}(M-k-\alpha)$ This symmetry condiction, along with the symmetry condictions for h(h), can be used to reduce the frequency specificallions from M- points to M+1 points for Model and 14/2 points for Meven.

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Unit sample response
$$k(w) = \pm k(M-1-w)$$

 $\frac{(4)}{(1+w)} = G_{T}(w) = \int_{T}^{T} k/M \quad k = 0,1,...,M-1$
 $G(w) = G_{T}(w) = \int_{T}^{T} k/M \quad k = 0,1,...,M-1$
 $G(w) = (-1)^{k} H_{T} \left(\frac{2\pi k}{H}\right) \quad G(w) = -G(H-4)$
 $k(w) = \frac{1}{H} \int_{V}^{1} G(w) + \frac{1}{2} \int_{K_{m}}^{\infty} G(w) \cos \frac{2\pi k}{H} (m+\frac{1}{2}) \int_{V}^{1}$
 $U = \int_{U}^{1} \frac{H}{H} \int_{U}^{1} \dots M - \sigma dd$
 $H(w + \frac{1}{2}) = G(w + \frac{1}{2}) = \int_{T}^{1} \frac{1}{2} \frac{\pi}{\pi} (k+\frac{1}{2}) \int_{U}^{1}$
 $w = \frac{1}{M} \quad G(w + \frac{1}{2}) = G(w + w + \frac{1}{2})$
 $k(w) = \frac{1}{M} \int_{K_{m}}^{\infty} G(w + \frac{1}{2}) s_{N} \frac{2\pi}{M} (k+\frac{1}{2}) (n+\frac{1}{2})$
 $k(w) = \frac{1}{M} \int_{K_{m}}^{\infty} G(w + \frac{1}{2}) s_{N} \frac{2\pi}{M} (k+\frac{1}{2}) (n+\frac{1}{2})$
 $H(w) = G(w) = \int_{T}^{1} \frac{\pi}{2} \int_{T}^{1} \frac{1}{2} \frac{1}{M} \frac{1}{M} (w + \frac{1}{2}) (n+\frac{1}{2})$
 $k(w) = (-1)^{k} H_{T} \left(\frac{2\pi k}{M}\right),$
 $k(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{K_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$
 $h(w) = -\frac{1}{M} \int_{W_{m}}^{\infty} G(w) s_{N} \frac{2\pi k}{M} (n+\frac{1}{2}), m - \omega dd$

$$\begin{aligned} & \text{EnggTree.com}_{H(k+\frac{1}{2})=a(k+\frac{1}{2})=J^{\pi}(2k+D/2H} \\ & \text{w}=\frac{1}{2} \quad Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=Q(k+\frac{1}{2})=$$

symmetric unit suite fies the conditions
response that settingies the conditions
$$k=0,1,2,3$$

 $H_r\left(\frac{2\pi k}{1\pi}\right) = \int_{0.4}^{1} k=4$
 $k=5,6,7$

Sodulios:
Since him & symmetric and frequencies are
selected to correspond to the case
$$\alpha = 0$$
, use the
symmetric formula to find her).
 $G(le) = (-1)^{le} H_0 \left(\frac{2\pi le}{H_1}\right)$
 $G(le) = 1$ $G(l) = -1$ $G(l) = 1$ $G(l) = -1$
 $G(l) = 1$ $G(l) = -1$ $G(l) = 0$ $G(l) = 0$
 $G(h) = 0.4$ $G(h) = 0$ $G(l) = 0$ $G(l) = 0$
 $h(m) = \frac{1}{15} \int_{-15}^{15} G(lo) + 2 \leq C G(le) \cos \left[\frac{2\pi le}{H} C n + \frac{1}{2}\right] \int_{-15}^{15} C G(l) + 2 \leq C G(le) \cos \left[\frac{2\pi le}{H} C n + \frac{1}{2}\right] \int_{-15}^{15} C G(l) + 2 \leq C G(le) \cos \left[\frac{2\pi le}{H} C n + \frac{1}{2}\right] \int_{-15}^{15} C G(l) + 2 = 7$

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$$h(n) = \frac{1}{15} \begin{cases} 1 + 2 & a(le) & cos \left[\frac{2\pi k}{15} (n + \frac{1}{2}) \right] \end{cases}$$

$$h(n) = \frac{1}{15} \begin{cases} 1 + 2 \left[(-1) (0 \cdot 7781) + 1 (0 \cdot 91354) + (-1) (0 \cdot 809) + (0 \cdot 669) \right] \end{cases}$$

$$= \frac{1}{15} \begin{cases} -0 \cdot 21(92) + (-1) (0 \cdot 809) + (0 \cdot 669) + (-1) (0 \cdot 669) + (-1) (0 \cdot 669) \end{bmatrix}$$

$$= \frac{1}{15} \begin{cases} -0 \cdot 21(92) + (-1) (0 \cdot 809) + (0 \cdot 669) + (-1) (0 \cdot 669) \end{bmatrix}$$

$$h(n) = -0 \cdot 0 \cdot 019 + 5309 = h(13)$$

$$h(n) = -0 \cdot 0019 + 5309 = h(13)$$

$$h(2) = 0 \cdot 019 + 5309 = h(13)$$

$$h(2) = 0 \cdot 019 + 5309 = h(13)$$

$$h(3) = 0 \cdot 012 + 34554 = h(11)$$

$$h(4) = -0 \cdot 09138802 = h(16)$$

$$h(5) = -0 \cdot 01808986 = h(9)$$

$$h(6) = 0 \cdot 3133176 = h(18)$$

$$h(7) = 0 \cdot 52$$

$$H\left(\frac{2\pi k}{15}\right) = 1$$
 for $k = 0, 1, 2, 3$
 0 for $k = 4, 5, 6, 7$ (A) $[m - m]$

Solution $H\left(\frac{2\pi k}{1\pi}\right) = \int I \qquad k = 0, 1, 2, 3 \qquad K = 0$ $Zo \qquad k = 4, 5, 6, 7$ $G(k) = (-1)^{k} H\left(\frac{2\pi k}{15}\right)$

$$C_{1}(0) = \underbrace{\text{EnggTree.com}}_{(a,b) = -1}^{(a,b) = -1}_{(a,b) = -1}^{(a,b) = -1}_{(b,b) = -1}$$

EnggTree.com FOURIER SERIES HETHOD The sinusoidal steady-state transfer functions of a digital filter & poriodic in the sampling frequency and if can be expanded in a Fourier serves. Thus $H(z)/z = j\omega = H(z)\omega$ = 5 hense _ 0 - - 00 where h(n) represents the terms of the unit impose response. In the equation O the filter is assumed to be non-causal and it can be modified to gread a causal filter. Rearranging the transfer function given below into real and imaginary components, we get $H(e^{j\omega}) = \frac{\infty}{5}$ here coscont $-j = \frac{\infty}{n-\infty}$ here $\omega_n \tau$ $H(a)(a) = H_r(4) - jH_r(4)$ lè, Hrlep = & Llos cos 2 #for _@ where and Hilt = 3 how sin 2 # for 3 Hr (4) & an even function and Hilf) & an odd function of frequency. If hent) is an even sequence, the inaginary part of the transfor function, Hilf) be zero and if h(nT) to an odd sequence, the real part of the transfer fundios, Holy will be zero. will Thus an over unit impulse response yields a real transfer function and ar odd unet impulse response yields an smag snarg fransfer feindion.

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DESIGN EQUATIONS

The storm H (el") is periodic in the sampling frequency and hence both Hrof and Hild are also periodic in the sampling frequency. Since the real part of the TF, Hr (4) is an even fanctions of frequency, its Fourier serve will be of atte form at Hr(f)= ao f & an coo(2 Tfn T) h=1 The Fourier co-eff. an are given by file $an = \frac{2}{fs} \int H_s (4) \cos\left(2\pi 4nT\right) df, n \neq 0$ -4s/2 and the as term is given by $a_0 = \frac{1}{f_s} \int \frac{1}{F_s} \frac{1}{f_$ Simillarly, the imaginary part of the TF, which Q an odd fanction of frequency can be expanded in the Fourier serves H+(+)= 5 5n STN (2 F. 4 AT) The Fourier coeff. In are given by Blz $b_n = \frac{2}{f_s} \int H_s(f) sin(2\pi f_n \tau) df$ $\frac{b_{0}=0}{For Hildson} = a_{0} + \frac{5}{n_{1}} \frac{a_{n}}{2} \left(2^{n} + 2^{-n}\right) / 2 = e^{\frac{1}{2}r f T}$ The TF can also be coritien as, $H(e^{j\omega}) = k(0) + \frac{5}{n-1} \left[h(-n) = \frac{n}{n+1} + k(n) = \frac{n}{n-1} \right]$ hlo) = 20 $h(-n) = \frac{1}{2} a_n \frac{1}{2} n = \frac{1}{2} a_n$

EnggTree.com For Hr(4)20 $H(e^{j_2}ff) = \frac{5}{n-1} \frac{b_n}{2} \left[\frac{2^n - 2^{-n}}{2} \right] / \frac{1}{2} e^{j_2}ff$ $h(n) = \frac{1}{2} bn \int n x o$ $h(n) = -\frac{1}{2} bn \int$ (Decide whether Hr(4) con High & to be Design Procedure: set equal to zero. For filtering applications we typically set Hilf =0. € Expand Hr(\$) or Hi(\$) In a Fourier Servis 3 The unif impulse response to determined from the Fourcer co-efficients using equal $H_{1}^{i}(\underline{t}) = 0 \qquad h(-K) = \frac{1}{2} a_{n} \qquad h(-K)$ $H_{r}(4) \equiv 0$ $h(-n) = \frac{1}{2} bn \int h > 0$ $h(n) = -\frac{1}{2} bn \int h > 0$ O Use the Fourier servis method to design a does-pa digital filter to approximate the ideal specification given by $H(e^{j\omega}) = \int_{0}^{1} \frac{4\sigma}{4r} \frac{1415fp}{4rc}$ where fp is the passband frequency and F is the sampling froquency.







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Lattice Structure:
The H-1

$$y(n) = x(n) + x_1(n) x(n-1)$$

 $k_1 = \infty_1(n) \rightarrow Reflection co-affroant.$
 $\downarrow (n) = x_1(n) + Reflection co-affroant.$
 $\downarrow (n) = y_1(n) = x_1(n) + y_1(n) = y_1(n)$
 $\downarrow (n) = \frac{1}{y_1(n)} + \frac{1}{y_1(n)} = x_1(n) + \frac{1}{x_1} x_1(n-1)$
 $g_1(n) = \frac{1}{x_1} + \frac{1}{y_2(n)} = x_1(n) + \frac{1}{x_1} x_1(n) + \frac{1}{x_1(n-1)}$
 $d_n general.$
 $\oint (n) = \frac{1}{y_1(n)} + \frac{1}{y_1(n)} + \frac{1}{y_1(n-1)} = \frac{1}{y_1(n)} + \frac{1}{y_1(n)}$



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Poly Phase Realization of Fie (1000)

$$H(z) = \sum_{n=0}^{m-1} k(n) z^{-n}$$

 $H(z) = \frac{m-1}{n>0}$
 $H(z) = \frac{m-1}{n>0}$
 $H(z) = h(0) z^{-1} + h(z) z^{-1} + h(z) z^{-4} + h(z) z^{-6} + h(z) z^{-6}$
 $+ h(0) z^{-1} + h(z) z^{-7} + h(z) z^{-6} + h(z) z^{-6} + h(z) z^{-6}$
 $Divide + ka i h = z subsystems
 $H(z) = h_0 (z^2) + z^{-1} P_1(z^2)$
 $h_0(z) = h(0) + h(z) z^{-1} + h(z) z^{-4} + h(z) z^{-6} + h(z) z^{-6}$
 $P_0(z) = h(0) + h(z) z^{-1} + h(z) z^{-4} + h(z) z^{-3} + h(z) z^{-6}$
 $H(z) = \sum_{m=0}^{m-1} z^{-m} h_m(z^m)$
 $wRue = \frac{m+1}{P_m(z^m)} + (Mh+m) z^m = 0 \le m \le m-1$
 $P_m(z^m) = \sum_{m=0}^{m-1} h(Mh+m) z^m = 0 \le m \le m-1$
 $\frac{1}{P_m(z^m)} + \frac{1}{P_0(z^m)} + \frac{1}{2} + \frac{1$$

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UNIT – IV FINITE WORD LENGTH EFFECTS

- Represent the following numbers in floating point format with five bits in mantissa and three bits in exponent.
 (8) [M/J – 13 R08]
 - a) 7₁₀
 - **b**) 0.25₁₀
 - c) -1₁₀
 - d) -0.25_{10}

Solution:

a) 7₁₀

 $(7)_{10} = 0.875 \text{ x } 2^3 = 0.1110 \text{ x } 2^{011}$

b) 0.25₁₀

 $(0.25)_{10} = 0.125 \text{ x } 2^1 = 0.0010 \text{ x } 2^{001}$

c) -1₁₀

Assume sign magnitude representation for negative number

 $(-1)_{10} = -0.5 \text{ x } 2^1 = 1.1000 \text{ x } 2^{001}$

d) -0.25_{10}

Assume sign magnitude representation for negative number

$$(-0.25)_{10} = -0.125 \text{ x } 2^1 = 1.0010 \text{ x } 2^{001}$$

2. Distinguish between fixed point and floating point arithmetic. (4)

[N/D - 11 R08] [M/J - 12 R08]

S.No.	Fixed Point Arithmetic	Floating Point Arithmetic		
1	Fast Operation	Slow Operation		
2	Relatively economical	More expensive because of costlier hardware		
3	Small dynamic range	Increased dynamic range		
4	Round off error occur only for additions	Round off errors can occur with both additions and multiplication		
5	Overflow occur in addition	Overflow does not arise		
6	Used in small computers	Used in general purpose computers		

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3. Discuss the various common methods of quantization. (8) [1

(8) [N/D – 13 R08]

The common methods of quantization are

- i. Truncation
- ii. Rounding

Truncation:

Truncation is process of reducing the size of binary number by discarding all bits less significant than the least significant bit that is retained. In the truncation of a binary number to b bits, all the less significant bits beyond bth bit are discarded.

Examples:

- i. The 8 bit binary number 0.00110011 may be truncated to 4 bits as 0.0011
- ii. The 8 bit binary number 1.01001001 may be truncated to 4 bits as 1.0100

When a number is truncated, the signal value is approximated by the highest quantization level that is not greater than the signal.

Rounding:

Rounding is the process of reducing the size of a binary number to finite word size of b bits such that, the rounded b bit number is closest to the original unquantized number.

Examples:

- i. The binary number 0.11010 may be rounded to three bits as either 0.110 or 0.111.
- ii. If the binary number 0.110111111 is rounded to 8 bits then the result may be 0.11011111 or 0.1110000.

Rounding up or down will have negligible effect on accuracy of computation.

4. Explain the errors due to rounding and truncation. [N/D - 10 R08] Explain the problems due to round off and truncation in converting a decimal fraction. (8) [N/D - 10 R08] Compare the truncation and rounding errors using fixed point and floating point representation. (8) [M/J - 14 R08]

Rounding or truncation introduces an error whose magnitude depends on the number of bits truncated or rounded-off. Also, the characteristics of the error depend on the form of binary number representation.

Consider a number x, whose original length is 'L' bits, as shown below.

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Let this number be quantized (truncated or rounded) to 'B' bits, as shown below.



The quantized number is represented as Q(x). Here B < L.

i) Truncation error for sign magnitude representation:

When the number x is positive, truncation results in reducing the magnitude of the number. Thus the truncation error is negative and the range is given by,

$$-(2^{-B} - 2^{-L}) \le \varepsilon_T \le 0 \tag{1}$$

The largest error occurs when all the discarded bits are one's.

When the number x is negative, truncation results in reducing the magnitude of the number. Because of the negative number the resulting will be greater than the original number. Thus the truncation error is positive and the range is given by,

$$0 \le \varepsilon_T \le (2^{-B} - 2^{-L}) \tag{2}$$

The overall range for the sign magnitude representation is

$$-(2^{-B} - 2^{-L}) \le \varepsilon_T \le (2^{-B} - 2^{-L})$$
(3)

ii) Truncation error for two's complement representation:

When the number x is positive, truncation results in reducing the magnitude of the number as in the case of sign magnitude numbers. Thus the truncation error is negative. When the number x is negative, truncation results in smaller number. Thus the truncation error is negative. The complete range of truncation error for two's complement is given by,

$$-(2^{-B} - 2^{-L}) \le \varepsilon_T \le 0 \tag{4}$$

iii) Round-off error for sign magnitude and two's complement representation:

The rounding of a binary number involves only the magnitude of the number and is independent of the type of fixed point binary representation. The error due to rounding may be either positive or negative. Its range is

$$-\frac{(2^{-B}-2^{-L})}{2} \le \varepsilon_R \le \frac{(2^{-B}-2^{-L})}{2}$$
(5)

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In most of the applications, infinite precision is assumed i.e., $L = \infty$

Therefore,

i) Truncation error for sign magnitude representation

 $-2^{-B} \le \epsilon_T \le 2^{-B}$

ii) Truncation error for two's complement representation

 $-2^{-B} \le \epsilon_T \le 0$

iii) Rounding error for sign magnitude and two's complement representation

$$\frac{-2^{-B}}{2} \le \epsilon_T \le \frac{2^{-B}}{2}$$

5. What is quantization noise? Derive the expression for quantization noise power. (12) [M/J - 12 R08]

What is meant by quantization? Derive the expression for the quantization error.

[N/D - 12 R08]Explain the quantization noise and derive the expression for finding quantization noise power. (8) [N/D - 10 R08]

QUANTIZATION:

In DSP, the continuous time input signals are converted into digital using ADC. The process of converting analog signal to digital signal is given below.



The signal x(t) is sampled at regular intervals t = nT, where n = 0, 1, 2, ... to create a sequence x(n). This is done by a sampler. The numeric equivalent of each sample x(n) is expressed by a finite number of bits giving the sequence $x_a(n)$. The

difference signal $e(n) = x_q(n) - x(t)$ is called quantization noise or A/D conversion noise.

The common methods of quantization are

- 1) Truncation
- 2) Rounding

EXPRESSION FOR QUANTIZATION NOISE POWER:

Input quantization error:

The quantization error is given by

$$e(n = x_0 n - x_0)$$

where

 $x_q(n)$ = sampled quantization value

x(n) = sampled unquantized value

Depending on the way in which (x) is quantized different distributions of quantization noise may be obtained. If rounding is used, the error signal satisfies the relation

$$\frac{-2^{-B}}{2} \le e(n) \le \frac{2^{-B}}{2}$$

The other type of quantization can be obtained by truncation. In truncation the signal is represented by the highest level that is not greater than the signal. In two's complement truncation, the error \mathbf{e}) is always negative and satisfies the inequality,

 $-2^{-B} \le e(n < 0)$

Steady state input noise power:

The quantization error is commonly viewed as an additive noise signal, that is

$$x_q(n = k) n + (e n)$$

The A/D converter output is the sum of the input signal x(n) and the error signal e(n) as shown below.



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If rounding is used for quantization error $e(n) = x_q(n) - x(n)$ is bounded by $\frac{-2^{-B}}{2} \le e(n) \le \frac{2^{-B}}{2}$.

The error e(n) has the following properties:

- 1. The error sequence e(n) is a sample sequence of a stationary random process.
- 2. The error sequence is uncorrelated with x(n) and other signals in the system.
- 3. The error is a white noise process with uniform amplitude probability distribution over the range of quantization error.

The variance of e(n) is given by

$$\sigma_e^2 = E[e^2(n)] - E^2[e(n)]$$

where $E^2[e(n)]$ is the average of $e^2(n)$ and E[e(n)] is mean value of e(n)

For rounding

$$\sigma_e^2 = \int_{\frac{-2^{-B}}{2}}^{\frac{2^{-B}}{2}} e^2 p_e(n) de - (0)^2$$
$$\sigma_e^2 = 2^B \left[\frac{e^3}{3}\right]_{\frac{-2^{-B}}{2}}^{\frac{2^{-B}}{2}}$$

Therefore,

$$\sigma_e^2 = \frac{2^{-2B}}{12}$$

6. Derive the signal to quantization noise ratio of A/D converter. (6) [M/J – 14 R08]

The signal to quantization noise ratio of A/D converter is given by

$$SNR = 10 \log[Px(n)/Pe(n)]$$

= 10 log Px(n) - 10 log Pe(n)
= 10 log Px(n) - 10 log $\frac{2^{-2B}}{12}$
= 10 log Px(n) + 10 log 2^{-2B} + 10 log 12
= 10 log Px(n) + 20B log 2+ 10.8
= 10 log Px(n) + 6.02B + 10.8
SNR = 10 log Px(n) + 6B + 10.8

For every increase in B bits the signal to quantization noise ratio increases by 6 dB.

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7. How is the steady state output noise variance calculated? (8) [N/D – 10 R08]

Steady state output noise variance (or) power:

The quantized input to a digital system with impulse response h(n) due to A/D conversion noise can be represented as shown in figure.



Let $\varepsilon(n)$ be the output noise due to quantization of the input. Then,

$$\varepsilon(n) = e(n) * h(n)$$
$$= \sum_{k=0}^{n} h(k)e(n-k)$$

The variance of any term in above sum in equal to $\sigma^2 h^2(n)$.

The variance of the sum of independent random variable is the sum of their variances. If the quantization errors are assumed to be independent at different sampling instances, then the variance of the output is

$$\sigma_{\varepsilon}^{2}(n)(or)\sigma_{e0}^{2}(n) = \sigma_{e}^{2}\sum_{n=0}^{k}h^{2}(n)$$

To find the steady state variance, extend the limit k up to infinity. Then,

$$\sigma_{\varepsilon}^{2} = \sigma_{e}^{2} \sum_{n=0}^{\infty} h^{2}(n)$$

Using Parseval's theorem the steady state output noise variance due to the quantization error is given by

$$\sigma_{\varepsilon}^{2} = \sigma_{e}^{2} \sum_{n=0}^{\infty} h^{2}(n) = \sigma_{e}^{2} \frac{1}{2\pi j} \oint_{c} H(z) H(z^{-1}) z^{-1} dz$$

where the closed contour of integration is around the unit circle |z| = 1 in which case only the poles that lie inside the unit circle are evaluated using the residue theorem.

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8. Determine the steady state output noise variance due to quantization of input, for the first order filter y(n) = ay(n-1) + x(n). (16)

Given y(n) = ay(n-1) + x(n)

Taking z transform on both sides

$$\begin{aligned} y(z) &= az^{-1}y(z) + x(z) \\ H(z) &= \frac{y(z)}{x(z)} = \frac{1}{1 - az^{-1}} = \frac{z}{z - a} \\ H(z^{-1}) &= \frac{z^{-1}}{z^{-1} - a} \\ \sigma_{\epsilon}^{2} &= \sigma_{\epsilon}^{2} \frac{1}{2\pi j} \oint_{c} H(z)H(z^{-1}) z^{-1} dz \\ \sigma_{\epsilon}^{2} &= \sigma_{\epsilon}^{2} \frac{1}{2\pi j} \oint_{c} \frac{z}{z - a} \frac{z^{-1}}{z^{-1} - a} z^{-1} dz \\ \sigma_{\epsilon}^{2} &= \sigma_{\epsilon}^{2} \frac{1}{2\pi j} \oint_{c} \frac{z^{-1}}{(z - a)(z^{-1} - a)} dz \\ \sigma_{\epsilon}^{2} &= \sigma_{\epsilon}^{2} \left[residue \ of \ \frac{z^{-1}}{(z - a)(z^{-1} - a)} \ at \ (z = a) + residue \ of \ \frac{z^{-1}}{(z - a)(z^{-1} - a)} \ at \ (z = \frac{1}{a}) \right] \\ \text{Assume } a < 1 \\ \sigma_{\epsilon}^{2} &= \sigma_{\epsilon}^{2} \left[\left((z - a) \frac{z^{-1}}{(z - a)(z^{-1} - a)} \right) \right]_{z = a} \\ \sigma_{\epsilon}^{2} &= \sigma_{\epsilon}^{2} \frac{a^{-1}}{a^{-1} - a} = \sigma_{\epsilon}^{2} \frac{1}{1 - a^{2}} \end{aligned}$$

9. Consider a second order IIR filter with $H(z) = \frac{1.0}{(1 - 0.5 z^{-1})(1 - 0.45 z^{-1})}$. Explain the effect of quantization on pole locations of the system when realized in direct form and in cascade form. Assume b = 3 bits. (10) [N/D - 11 R08]

Solution:

Given:
$$H(z) = \frac{1.0}{(1-0.5 \text{ z}^{-1})(1-0.45 \text{ z}^{-1})}$$

Direct form:

H(z) can be written as,

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$$H(z) = \frac{1}{(1 - 0.95z^{-1} + 0.225z^{-2})}$$

The poles are at $p_1 = 0.5$ and $p_2 = 0.45$

Let us quantize the co-efficient by truncation.

$$(0.95)_{10} \xrightarrow{\text{to binary}} (0.1111)_2 \xrightarrow{\text{truncate to 3 bits}} (0.111)_2 \xrightarrow{\text{to decimal}} (0.875)_{10}$$
$$(0.225)_{10} \xrightarrow{\text{to binary}} (0.0011)_2 \xrightarrow{\text{truncate to 3 bits}} (0.001)_2 \xrightarrow{\text{to decimal}} (0.125)_{10}$$

Hence,

$$H(z) = \frac{1}{(1 - 0.875z^{-1} + 0.125z^{-2})}$$
$$H(z) = \frac{z^2}{z^2 - 0.875z + 0.125}$$
$$H(z) = \frac{z^2}{(z - 0.695)(z - 0.18)}$$
So, $\overline{p_1} = 0.695$ and $\overline{p_2} = 0.18$

Thus the pole locations are shifted from 0.5, 0.45 to 0.695 and 0.18 respectively.

Cascade Form:

$$H(z) = \frac{1.0}{(1 - 0.5 z^{-1})(1 - 0.45 z^{-1})}$$

(0.5)₁₀ $\xrightarrow{\text{to binary}}$ (0.1000)₂ $\xrightarrow{\text{truncate to 3 bits}}$ (0.100)₂ $\xrightarrow{\text{to decimal}}$ (0.5)₁₀
(0.45)₁₀ $\xrightarrow{\text{to binary}}$ (0.0111)₂ $\xrightarrow{\text{truncate to 3 bits}}$ (0.011)₂ $\xrightarrow{\text{to decimal}}$ (0.375)₁₀

Hence,

$$H(z) = \frac{1.0}{(1 - 0.5 z^{-1})(1 - 0.375 z^{-1})}$$

So, $\overline{p_1} = 0.5$ and $\overline{p_2} = 0.375$

There is no change in first pole location and the second pole location is shifted from 0.45 to 0.375.

10. Explain coefficient quantization in IIR filter. (16) [N/D – 12 R08]

The filter coefficients are computed to infinite precision in theory. But, in digital computation the filter coefficients are represented in binary and are stored in registers. If b bit register is used, the filter coefficients must be rounded or truncated to b bits, which produces an error. This error is known as coefficient quantization error.

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Due to quantization of coefficients, the frequency response of the filter may differ appreciably from the desired response and sometimes the filter may actually fail to meet the desired specifications. If the poles of desired filter are close to the unit circle, then those of the filter with quantized coefficients may lie just outside the unit circle, leading to instability.

Example:

Consider a second order IIR filter with $H(z) = \frac{1.0}{(1 - 0.5 z^{-1})(1 - 0.45 z^{-1})}$. Explain the effect of quantization on pole locations of the system when realized in direct form and in cascade form. Assume b = 3 bits.

Answer: Refer problem No: 9

11. Draw the product quantization noise model of second order IIR system. (8)

 $\left[M/J-13\;R08\right]$



Quantization noise model for a second-order system with five noise



Quantization noise model for a second order system with a single noise

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12. Determine the output round-off noise power for the system having transfer function $H(z) = \frac{1}{(1-0.5z^{-1})(1-0.4z^{-1})}$ which is realized in cascade form. Assume word

length is 4 bits. (8)

Solution:

Given:



$$\begin{aligned} \sigma_{e0}^2 &= \sigma_{e01}^2 + \sigma_{e02}^2 \\ \sigma_{e01}^2 &= \sigma_e^2 \frac{1}{2\pi j} \oint_c H(z) H(z^{-1}) \, z^{-1} \, dz \\ \sigma_{e01}^2 &= \sigma_e^2 I \end{aligned}$$

where

$$I = \frac{1}{2\pi j} \oint_{c} H(z) H(z^{-1}) z^{-1} dz$$

H(z) is the transfer function seen by the error source $e_1(n)$

$$H(z) = \frac{1}{(1 - 0.5z^{-1})(1 - 0.4z^{-1})}$$

I = sum of residues at poles within the unit circle

$$I = \frac{1}{2\pi j} \oint_{c} \frac{1}{(1 - 0.5z^{-1})(1 - 0.4z^{-1})} \frac{1}{(1 - 0.5z)(1 - 0.4z)} z^{-1} dz$$
$$I = \frac{1}{2\pi j} \oint_{c} \frac{(z)(z)}{(z - 0.5)(z - 0.4)} \frac{1}{(1 - 0.5z)(1 - 0.4z)} z^{-1} dz$$
$$I = \frac{1}{2\pi j} \oint_{c} \frac{(z)}{(z - 0.5)(z - 0.4)} \frac{1}{(1 - 0.5z)(1 - 0.4z)} dz$$

The poles z = 0.5 and z = 0.4 are within unit circle.

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$$I = (z - 0.5) \frac{(z)}{(z - 0.5)(z - 0.4)} \frac{1}{(1 - 0.5z)(1 - 0.4z)}\Big|_{z=0.5} + (z - 0.4) \frac{(z)}{(z - 0.5)(z - 0.4)} \frac{1}{(1 - 0.5z)(1 - 0.4z)}\Big|_{z=0.4}$$

$$I = 8.33 - 5.95 = 2.38$$

$$\sigma_{e01}^2 = \sigma_e^2 (2.38)$$

$$\sigma_{e01}^2 = \frac{2^{-2B}}{12} (2.38)$$

$$\sigma_{e01}^2 = \frac{2^{-2(4)}}{12} (2.38)$$

$$\sigma_{e01}^2 = 3.25X 10^{-4} (2.38)$$

$$\sigma_{e01}^2 = 7.747X 10^{-4}$$

$$\sigma_{e02}^2 = \sigma_e^2 \frac{1}{2\pi j} \oint_c H(z) H(z^{-1}) z^{-1} dz$$

$$\sigma_{e02}^2 = \sigma_e^2 I$$

where

$$I = \frac{1}{2\pi j} \oint_{c} H(z) H(z^{-1}) \, z^{-1} \, dz$$

H(z) is the transfer function seen by the error source $e_2(n)$

$$H(z) = \frac{1}{(1 - 0.4z^{-1})}$$

I = sum of residues at poles within the unit circle

$$I = \frac{1}{2\pi j} \oint_{c} \frac{1}{(1 - 0.4z^{-1})} \frac{1}{(1 - 0.4z)} z^{-1} dz$$
$$I = \frac{1}{2\pi j} \oint_{c} \frac{\frac{z}{(z - 0.4)}}{(z - 0.4)} \frac{1}{(1 - 0.4z)} z^{-1} dz$$
$$I = \frac{1}{2\pi j} \oint_{c} \frac{1}{(z - 0.4)} \frac{1}{(1 - 0.4z)} dz$$

The pole z = 0.4 is within unit circle.

$$I = (z - 0.4) \frac{1}{(z - 0.4)} \frac{1}{(1 - 0.4z)} \Big|_{z = 0.4}$$
$$I = 1.1905$$

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$$\sigma_{e02}^2 = \sigma_e^2 (1.1905)$$

$$\sigma_{e02}^2 = 3.25X10^{-4} (1.1905)$$

$$\sigma_{e02}^2 = 3.869X10^{-4}$$

Therefore,

$$\sigma_{e0}^2 = 7.747X10^{-4} + 3.869X10^{-4} = 11.616X10^{-4}$$

13. Determine the dead band of the system y(n) = 0.2y(n-1) + 0.5y(n-2) + x(n). Assume 8 bits are used for signal representation. (8) [M/J - 13 R08]

Given:
$$y(n) = 0.2y(n-1) + 0.5y(n-2) + x(n)$$

Assume

$$x(n) = \begin{cases} 0.0125, & n = 0 \\ 0, & otherwise \end{cases}$$
$$y(-1) = y(-2) = 0$$

b = 8 bits (excluding sign bit)

n	<i>x</i> (<i>n</i>)	0.2y(n-1)	Q[0.2y(n-1)]	0.5y(n-2)	Q[0.5y(n-2)]	y (n)
0	0.0125	0	0	0	0	0.0125
1	0	0.0025	0.00390625	0	0	0.00390625
2	0	0.00078125	0	0.00625	0.0078125	0.0078125
3	0	0.0015625	0	0.001953125	0.00390625	0.00390625
4	0	0.00078125	0	0.00390625	0.00390625	0.00390625

$$(0.0025)_{10} \xrightarrow{\text{to binary}} (0.000000010)_{2} \xrightarrow{\text{round off to 8 bits}} (0.00000001)_{2} \xrightarrow{\text{to decimal}} (0.00390625)_{10}$$

$$(0.00078125)_{10} \xrightarrow{\text{to binary}} (0.00000000)_{2} \xrightarrow{\text{round off to 8 bits}} (0.00000000)_{2} \xrightarrow{\text{to decimal}} (0)_{10}$$

$$(0.00625)_{10} \xrightarrow{\text{to binary}} (0.000000011)_{2} \xrightarrow{\text{round off to 8 bits}} (0.00000010)_{2} \xrightarrow{\text{to decimal}} (0.0078125)_{10}$$

$$(0.0015625)_{10} \xrightarrow{\text{to binary}} (0.00000000)_{10} + \xrightarrow{\text{round off to 8 bits}} (0.00000000)_{2} \xrightarrow{\text{to decimal}} (0)_{10}$$

$$0.001953125$$

$$\xrightarrow{\text{to binary}} (0.00000001)_{10} \xrightarrow{\text{round off to 8 bits}} (0.00000001)_{2} \xrightarrow{\text{to decimal}} (0.00390625)_{10}$$

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 $(0.00390625)_{10} \xrightarrow{\text{to binary}} (0.00000001)_{10} + \underbrace{\text{round off to 8 bits}}_{\text{round off to 8 bits}} (0.00000001)_2 \xrightarrow{\text{to decimal}} (0.00390625)_{10}$ Dead band $= \frac{\frac{1}{2}2^{-8}}{1-0.5}$

Dead band = 0.00390625

14. Describe the quantization in floating point realization of IIR digital filters.

(16) [N/D -13 R08]

(8) [M/J - 12 R08]

Refer Page No: 7.37 in Digital Signal Processing by Ramesh Babu

15. Explain the finite word length effects in FIR digital filters.(8)[N/D - 13 R08]Explain the round off noise in direct form realization of a linear phase FIR filterwith relevant diagrams.(8) [A/M - 11 R08]

Explain the effects of coefficient quantization in FIR filters.

(8)[M/J - 14 R08] [A/M -11 R08]

Refer Page No: 7.41 in Digital Signal Processing by Ramesh Babu

- 16. Explain the following
 - a) Coefficient quantization error
 - b) Product quantization error
 - c) Signal scaling
 - d) Truncation and Rounding

a) Coefficient quantization error:

The filter coefficients are computed to infinite precision in theory. But, in digital computation the filter coefficients are represented in binary and are stored in registers. If b bit register is used, the filter coefficients must be rounded or truncated to b bits, which produces an error. This error is known as coefficient quantization error. Due to quantization of coefficients, the frequency response of the filter may differ appreciably from the desired response and sometimes the filter may actually fail to meet the desired specifications. If the poles of desired filter are close to the unit circle, then those of the filter with quantized coefficients may lie just outside the unit circle, leading to instability.

b) Product quantization error:

Product quantization errors arise at the output of a multiplier. Multiplication of b bit data with b bit coefficient results in a product having 2b bits. Since b bit register

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is used, the multiplier output must be rounded or truncated to b bits, which produces an error. This error is known as product quantization error.

c) Signal scaling:

Saturation arithmetic eliminates limit cycles due to overflow, but it causes undesirable signal distortion due to the nonlinearity of the clipper. In order to limit the amount of non-linear distortion, it is important to scale the input signal (scaling factor S_0) and the unit sample response between the input and any internal summing nodes in the system such that overflow become a rare event.



d) Truncation and Rounding:

Truncation:

Truncation is process of reducing the size of binary number by discarding all bits less significant than the least significant bit that is retained. In the truncation of a binary number to b bits all the less significant bits beyond bth bit are discarded.

Rounding:

Rounding is the process of reducing the size of a binary number to finite word size of b bits such that, the rounded b bit number is closest to the original unquantized number.

17. a) How is signal scaling used to prevent overflow limit cycle in the digital filter implementation? Explain with an example. (8) (N/D 11 R08) (M/J 13 R08)

Signal Scaling:

Saturation arithmetic eliminates limit cycles due to overflow, but it causes undesirable signal distortion due to the nonlinearity of the clipper. In order to limit the amount of non-linear distortion, it is important to scale the input signal and the unit sample response between the input and any internal summing nodes in the system such that overflow become a rare event.

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Consider a second order IIR filter shown in figure. A scale factor S_0 is introduced between the input x(n) and the adder 1, to prevent overflow at the output adder 1.

Now the overall input-output transfer function is

$$H(z) = S_0 \frac{b_0 + b_1 z^{-1} + b_2 z^{-2}}{1 + a_1 z^{-1} + a_2 z^{-2}}$$
$$H(z) = S_0 \frac{N(z)}{D(z)}$$

From the figure,

$$H'(z) = \frac{W(z)}{X(z)} = \frac{S_0}{1 + a_1 z^{-1} + a_2 z^{-2}} = \frac{S_0}{D(z)}$$

If the instantaneous energy in the output sequence w(n) is less than the finite energy in the input sequence then, there will not be any overflow.

$$w(z) = \frac{S_0 X(z)}{D(z)} = S_0 S(z) X(z)$$

where,

$$S_{0}^{2} = \frac{1}{\frac{1}{2\pi j} \oint_{c} S(z) S(z^{-1}) z^{-1} dz}$$

$$S_{0}^{2} = \frac{1}{\frac{1}{2\pi j} \oint_{c} \frac{z^{-1} dz}{D(z) D(z^{-1})}}$$

$$S_{0}^{2} = \frac{1}{l}$$
where
$$I = \frac{1}{2\pi j} \oint_{c} \frac{z^{-1} dz}{D(z) D(z^{-1})}$$
18. Describe the effects of quantization in IIR filter. Consider a first order filter with difference equation y(n) = x(n) + 0.5y(n-1). Assume that the data register length is three bits plus a sign bit. The input $x(n)=0.875\delta(n)$. Explain the limit cycle oscillations in the above filter, if quantization is performed by means of rounding and signed magnitude representation. (16) Solution:

Given: y(n) = x(n) + 0.5y(n-1) $x(n) = 0.875 \,\delta(n)$ $n(n) = \begin{cases} 0.875, & n = 0 \\ 0 & otherwise \end{cases}$

$$x(n) = \{0, otherwise\}$$

$$y(-1)=0$$

b = 3 bits (excluding sign bit)

n	x(n)	0.5y(n-1)	Q[0.5y(n-1)]	y (n)
0	0.875	0	0	0.875
1	0	0.4375	0.5	0.5
2	0	0.25	0.25	0.25
3	0	0.125	0.125	0.125
4	0	0.0625	0.125	0.125

$$(0.4375)_{10} \xrightarrow{\text{to binary}} (0.0111)_2 \xrightarrow{\text{round off to 3 bits}} (0.100)_2 \xrightarrow{\text{to decimal}} (0.5)_{10}$$

$$(0.25)_{10} \xrightarrow{\text{to binary}} (0.0100)_2 \xrightarrow{\text{round off to 3 bits}} (0.010)_2 \xrightarrow{\text{to decimal}} (0.25)_{10}$$

$$(0.125)_{10} \xrightarrow{\text{to binary}} (0.0010)_2 \xrightarrow{\text{round off to 3 bits}} (0.001)_2 \xrightarrow{\text{to decimal}} (0.125)_{10}$$

$$(0.0625)_{10} \xrightarrow{\text{to binary}} (0.0001)_2 \xrightarrow{\text{round off to 3 bits}} (0.001)_2 \xrightarrow{\text{to decimal}} (0.125)_{10}$$

Dead band =
$$\frac{\frac{1}{2}2^{-3}}{1-0.5}$$

$$Dead \ band = 0.125$$

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19. Consider all pole second order IIR digital filter described by difference equation y(n) = -0.5y(n-1) - 0.75y(n-2) + x(n). Assuming 8-bits to represent a pole, determine the dead band region governing the limit cycle. (8) Given: y(n) = -0.5y(n-1) + 0.75y(n-2) + x(n)

Assume

$$x(n) = \begin{cases} 0.0125, & n = 0 \\ 0, & otherwise \end{cases}$$

$$y(-1) = y(-2) = 0$$

b = 8 bits (excluding sign bit)

n	x(n)	-0.5y(n-1)	Q[-0.5y(n-1)]	-0.75y(n-2)	Q[-0.75y(n	y (n)
					- 2)]	
0	0.0125	0	0	0	0	0.0125
1	0	0.0025	0.00390625	0	0	0.00390625
2	0	0.00078125	0	0.00625	0.0078125	0.0078125
3	0	0.0015625	0	0.001953125	0.00390625	0.00390625
4	0	0.00078125	0	0.00390625	0.00390625	0.00390625

$$(0.0025)_{10} \xrightarrow{\text{to binary}} (0.000000010)_2$$
$$\xrightarrow{\text{round off to 8 bits}} (0.00000001)_2 \xrightarrow{\text{to decimal}} (0.00390625)_{10}$$

$$(0.00078125)_{10} \xrightarrow{\text{to binary}} (0.00000000)_{2}$$
$$\xrightarrow{\text{round off to 8 bits}} (0.00000000)_{2} \xrightarrow{\text{to decimal}} (0)_{10}$$

 $(0.00625)_{10} \xrightarrow{\text{to binary}} (0.000000011)_2$ $\xrightarrow{\text{round off to 8 bits}} (0.00000010)_2 \xrightarrow{\text{to decimal}} (0.0078125)_{10}$

$$(0.0015625)_{10} \xrightarrow{\text{to binary}} (0.00000000)_{10} + \underbrace{\xrightarrow{\text{round off to 8 bits}}}_{0.000000000} (0.00000000)_2 \xrightarrow{\text{to decimal}} (0)_{10}$$

 $\xrightarrow{\text{to binary}} (0.00000001)_{10} \xrightarrow{\text{round off to 8 bits}} (0.00000001)_2 \xrightarrow{\text{to decimal}} (0.00390625)_{10}$

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 $(0.00390625)_{10} \xrightarrow{\text{to binary}} (0.00000001)_{10} + \underbrace{\text{round off to 8 bits}}_{\text{round off to 8 bits}} (0.00000001)_2 \xrightarrow{\text{to decimal}} (0.00390625)_{10}$ Dead band $= \frac{\frac{1}{2}2^{-8}}{1-0.75}$

 $Dead \ band = 0.0078125$

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EC 3492 – DIGITAL SIGNAL PROCESSING

UNIT V – DSP APPLICATIONS

- 1. A signal x(n) is given by $x(n) = \{0, 1, 2, 3, 4, 5, 6, 0, 1, 2, 3,\}$ (8)
 - a) Obtain the decimated signal with a factor of 2.
 - b) Obtain the interpolated signal with a factor of 2.

Solution:

Given: $x(n) = \{0, 1, 2, 3, 4, 5, 6, 0, 1, 2, 3, \dots\}$

a) Decimation with a factor of 2:

$$y(n) = \{0, 2, 4, 6, 1, 3, \dots\}$$

b) Interpolation with a factor of 2:

$$y(n) = \{0, 0, 1, 0, 2, 0, 3, 0, 4, 0, 5, 0, 6, 0, 0, 0, 1, 0, 2, 0, 3, 0, \dots\}$$

2. How does the sampling rate increase by an integer factor I and derive the input-

output relationship in both time and frequency domains.

An increase in the sampling rate by an integer factor of I can be accomplished by interpolating I - 1 new samples between successive values of the signal. The interpolation process can be accomplished in a variety of ways. We shall describe a process that preserves the spectral shape of the signal sequence x(n).

Let v(m) denote a sequence with a rate $F_y = IF_x$, which is obtained from x(n) by adding I - 1 zeros between successive values of x(n). Thus

$$v(m) = \begin{cases} x(m/I), & m = 0, \pm I, \pm 2I, \dots \\ 0, & \text{otherwise} \end{cases}$$
(10.3.1)

and its sampling rate is identical to the rate of y(m). This sequence has a z-transform

$$V(z) = \sum_{m=-\infty}^{\infty} v(m) z^{-m}$$

=
$$\sum_{m=-\infty}^{\infty} x(m) z^{-ml}$$

=
$$X(z^{l})$$
 (10.3.2)

The corresponding spectrum of v(m) is obtained by evaluating (10.3.2) on the unit circle. Thus

$$V(\omega_y) = X(\omega_y I) \tag{10.3.3}$$

where ω_y denotes the frequency variable relative to the new sampling rate F_y (i.e., $\omega_y = 2\pi F/F_y$). Now the relationship between sampling rates is $F_y = IF_x$ and hence, the frequency variables ω_x and ω_y are related according to the formula

$$\omega_{y} = \frac{\omega_{x}}{I} \tag{10.3.4}$$

The spectra $X(\omega_x)$ and $V(\omega_y)$ are illustrated in Fig. 10.5. We observe that the sampling rate increase, obtained by the addition of I - 1 zero samples between successive values of x(n), results in a signal whose spectrum $V(\omega_y)$ is an *I*-fold periodic repetition of the input signal spectrum $X(\omega_x)$.

Since only the frequency components of x(n) in the range $0 \le \omega_y \le \pi/I$ are unique, the images of $X(\omega)$ above $\omega_y = \pi/I$ should be rejected by passing the sequence v(m) through a lowpass filter with frequency response $H_I(\omega_y)$ that

ideally has the characteristic

$$H_I(\omega_y) = \begin{cases} C, & 0 \le |\omega_y| \le \pi/l \\ 0, & \text{otherwise} \end{cases}$$
(10.3.5)

where C is a scale factor required to properly normalize the output sequence y(m). Consequently, the output spectrum is





Figure 10.5 Spectra of x(n) and v(n) where $V(\omega_y) = X(\omega_y l)$.

The scale factor C is selected so that the output y(m) = x(m/l) for m = 0, $\pm l$, +2l,.... For mathematical convenience, we select the point m = 0. Thus

$$y(0) = \frac{1}{2\pi} \int_{-\pi}^{\pi} Y(\omega_y) d\omega_y$$

= $\frac{C}{2\pi} \int_{-\pi/I}^{\pi/I} X(\omega_y I) d\omega_y$ (10.3.7)

Since $\omega_y = \omega_x / I$, (10.3.7) can be expressed as

$$y(0) = \frac{C}{I} \frac{1}{2\pi} \int_{-\pi}^{\pi} X(\omega_x) d\omega_x$$

= $\frac{C}{I} x(0)$ (10.3.8)

Therefore, C = I is the desired normalization factor.

Finally, we indicate that the output sequence y(m) can be expressed as a convolution of the sequence v(n) with the unit sample response h(n) of the lowpass

filter. Thus

$$y(m) = \sum_{k=-\infty}^{\infty} h(m-k)v(k)$$
 (10.3.9)

Since v(k) = 0 except at multiples of *I*, where v(kI) = x(k), (10.3.9) becomes

$$y(m) = \sum_{k=-\infty}^{\infty} h(m - kI)x(k)$$
 (10.3.10)

3. Explain sampling rate reduction by an integer factor 'D'. Derive the relation between input and output frequency. (8)

DECIMATION BY A FACTOR D

Let us assume that the signal x(n) with spectrum $X(\omega)$ is to be downsampled by an integer factor D. The spectrum $X(\omega)$ is assumed to be nonzero in the frequency interval $0 \le |\omega| \le \pi$ or, equivalently, $|F| \le F_x/2$. We know that if we reduce the sampling rate simply by selecting every Dth value of x(n), the resulting signal will be an aliased version of x(n), with a folding frequency of $F_x/2D$. To avoid aliasing, we must first reduce the bandwidth of x(n) to $F_{\max} = F_x/2D$ or, equivalently, to $\omega_{\max} = \pi/D$. Then we may downsample by D and thus avoid aliasing.

The decimation process is illustrated in Fig. 10.2. The input sequence x(n) is passed through a lowpass filter, characterized by the impulse response h(n) and a frequency response $H_D(\omega)$, which ideally satisfies the condition

$$H_D(\omega) = \begin{cases} 1, & |\omega| \le \pi/D \\ 0, & \text{otherwise} \end{cases}$$
(10.2.1)

Thus the filter eliminates the spectrum of $X(\omega)$ in the range $\pi/D < \omega < \pi$. Of course, the implication is that only the frequency components of x(n) in the range $|\omega| \le \pi/D$ are of interest in further processing of the signal.

The output of the filter is a sequence v(n) given as

$$v(n) = \sum_{k=0}^{\infty} h(k)x(n-k)$$
(10.2.2)

x(n)	h(n)	<i>v(n)</i>	Downsampler ↓D	y(m) F.
$F_x = \overline{T_x}$				$F_y = \frac{1}{D}$

Figure 10.2 Decimation by a factor D.

which is then downsampled by the factor D to produce y(m). Thus

$$y(m) = v(mD)$$

= $\sum_{k=0}^{\infty} h(k)x(mD-k)$ (10.2.3)

Although the filtering operation on x(n) is linear and time invariant, the downsampling operation in combination with the filtering results in a time-variant system. This is easily verified. Given the fact that x(n) produces y(m), we note that $x(n-n_0)$ does not imply $y(n-n_0)$ unless n_0 is a multiple of D. Consequently, the overall linear operation (linear filtering followed by downsampling) on x(n) is not time invariant.

The frequency-domain characteristics of the output sequence y(m) can be obtained by relating the spectrum of y(m) to the spectrum of the input sequence x(n). First, it is convenient to define a sequence $\tilde{v}(n)$ as

$$\tilde{v}(n) = \begin{cases} v(n), & n = 0, \pm D, \pm 2D, \dots \\ 0, & \text{otherwise} \end{cases}$$
(10.2.4)

Clearly, $\tilde{v}(n)$ can be viewed as a sequence obtained by multiplying v(n) with a periodic train of impulses p(n), with period D, as illustrated in Fig. 10.3. The discrete Fourier series representation of p(n) is

$$p(n) = \frac{1}{D} \sum_{k=0}^{D-1} e^{j2\pi kn/D}$$
(10.2.5)

Hence

$$\tilde{v}(n) = v(n)p(n) \tag{10.2.6}$$

and

$$v(m) = \tilde{v}(mD) = v(mD)p(mD) = v(mD) \tag{10.2.7}$$



Figure 10.3 Multiplication of v(n) with a periodic impulse train p(n) with period D = 3.

Now the z-transform of the output sequence y(m) is

$$Y(z) = \sum_{m=-\infty}^{\infty} y(m) z^{-m}$$

=
$$\sum_{m=-\infty}^{\infty} \tilde{v}(mD) z^{-m}$$
 (10.2.8)
$$Y(z) = \sum_{m=-\infty}^{\infty} \tilde{v}(m) z^{-m/D}$$

where the last step follows from the fact that $\tilde{v}(m) = 0$, except at multiples of D. By making use of the relations in (10.2.5) and (10.2.6) in (10.2.8), we obtain

$$Y(z) = \sum_{m=-\infty}^{\infty} v(m) \left[\frac{1}{D} \sum_{k=0}^{D-1} e^{j2\pi mk/D} \right] z^{-m/D}$$

$$= \frac{1}{D} \sum_{k=0}^{D-1} \sum_{m=-\infty}^{\infty} v(m) (e^{-j2\pi k/D} z^{1/D})^{-m}$$

$$= \frac{1}{D} \sum_{k=0}^{D-1} V(e^{-j2\pi k/D} z^{1/D})$$

$$= \frac{1}{D} \sum_{k=0}^{D-1} H_D(e^{-j2\pi k/D} z^{1/D}) X(e^{-j2\pi k/D} z^{1/D})$$

(10.2.9)

where the last step follows from the fact that $V(z) = H_D(z)X(z)$.

By evaluating Y(z) in the unit circle, we obtain the spectrum of the output signal y(m). Since the rate of y(m) is $F_y = 1/T_y$, the frequency variable, which we denote as ω_y , is in radians and is relative to the sampling rate F_y ,

$$\omega_{y} = \frac{2\pi F}{F_{y}} = 2\pi F T_{y} \tag{10.2.10}$$

Since the sampling rates are related by the expression

$$F_y = \frac{F_x}{D} \tag{10.2.11}$$

it follows that the frequency variables ω_v and

$$\omega_x = \frac{2\pi F}{F_x} = 2\pi F T_x \tag{10.2.12}$$

are related by

$$\omega_y = D\omega_x \tag{10.2.13}$$

Thus, as expected, the frequency range $0 \le |\omega_x| \le \pi/D$ is stretched into the corresponding frequency range $0 \le |\omega_y| \le \pi$ by the downsampling process.

We conclude that the spectrum $Y(\omega_y)$, which is obtained by evaluating (10.2.9) on the unit circle, can be expressed as

$$Y(\omega_y) = \frac{1}{D} \sum_{k=0}^{D-1} H_D\left(\frac{\omega_y - 2\pi k}{D}\right) X\left(\frac{\omega_y - 2\pi k}{D}\right)$$
(10.2.14)

With a properly designed filter $H_D(\omega)$, the aliasing is eliminated and, consequently, all but the first term in (10.2.14) vanish. Hence



for $0 \le |\omega_y| \le \pi$. The spectra for the sequences x(n), v(n), and y(m) are illustrated in Fig. 10.4.

4. For the multirate system shown in figure, find the relation between x(n) and y(n).

(N/D 11 R08)



SOLUTION:



5. Explain sampling rate conversion by a rational factor and derive input-output relation in both time and frequency domain. (10)

SAMPLING RATE CONVERSION BY A RATIONAL FACTOR I/D

Having discussed the special cases of decimation (downsampling by a factor D) and interpolation (upsampling by a factor I), we now consider the general case of sampling rate conversion by a rational factor I/D. Basically, we can achieve this sampling rate conversion by first performing interpolation by the factor I and then decimating the output of the interpolator by the factor D. In other words, a sampling rate conversion by the rational factor I/D is accomplished by cascading an interpolator with a decimator, as illustrated in Fig. 10.6.

We emphasize that the importance of performing the interpolation first and the decimation second, is to preserve the desired spectral characteristics of x(n). Furthermore, with the cascade configuration illustrated in Fig. 10.6, the two filters with impulse response $\{h_u(l)\}$ and $\{h_d(l)\}$ are operated at the same rate, namely IF_x and hence can be combined into a single lowpass filter with impulse response h(l)as illustrated in Fig. 10.7. The frequency response $H(\omega_v)$ of the combined filter must incorporate the filtering operations for both interpolation and decimation, and hence it should ideally possess the frequency response characteristic

$$H(\omega_v) = \begin{cases} I, & 0 \le |\omega_v| \le \min(\pi/D, \pi/I) \\ 0, & \text{otherwise} \end{cases}$$
(10.4.1)

where $\omega_v = 2\pi F/F_v = 2\pi F/IF_x = \omega_x/I$.



Figure 10.6 Method for sampling rate conversion by a factor I/D.



Figure 10.7 Method for sampling rate conversion by a factor I/D.

In the time domain, the output of the upsampler is the sequence

$$v(l) = \begin{cases} x(l/I), & l = 0, \pm I, \pm 2I, \dots \\ 0, & \text{otherwise} \end{cases}$$
(10.4.2)

and the output of the linear time-invariant filter is

$$w(l) = \sum_{k=-\infty}^{\infty} h(l-k)v(k)$$

$$= \sum_{k=-\infty}^{\infty} h(l-kI)x(k)$$
(10.4.3)

Finally, the output of the sampling rate converter is the sequence $\{y(m)\}$, which is obtained by downsampling the sequence $\{w(l)\}$ by a factor of D. Thus

$$w(m) = w(mD)$$

= $\sum_{k=-\infty}^{\infty} h(mD - kI)x(k)$ (10.4.4)

It is illuminating to express (10.4.4) in a different form by making a change in variable. Let

$$k = \left\lfloor \frac{mD}{l} \right\rfloor - \frac{n}{l} \tag{10.4.5}$$

where the notation $\lfloor r \rfloor$ denotes the largest integer contained in r. With this change in variable, (10.4.4) becomes

$$y(m) = \sum_{n=-\infty}^{\infty} h\left(mD - \left\lfloor \frac{mD}{I} \right\rfloor I + nI\right) x\left(\left\lfloor \frac{mD}{I} \right\rfloor - n\right)$$
(10.4.6)

We note that

$$mD - \left\lfloor \frac{mD}{I} \right\rfloor I = mD \mod I$$
$$= (mD)_{I}$$

Consequently, (10.4.6) can be expressed as

$$y(m) = \sum_{n=-\infty}^{\infty} h(nI + (mD)_I) x\left(\left\lfloor \frac{mD}{I} \right\rfloor - n\right)$$
(10.4.7)

It is apparent from this form that the output y(m) is obtained by passing the input sequence x(n) through a time-variant filter with impulse response

$$g(n,m) = h(nI + (mD)_I) \qquad -\infty < m, n < \infty \tag{10.4.8}$$

where h(k) is the impulse response of the time-invariant lowpass filter operating at the sampling rate IF_x . We further observe, that for any integer k,

$$g(n, m + kI) = h(nI + (mD + kDI)_{I})$$

= $h(nI + (mD)_{I})$ (10.4.9)
= $g(n, m)$

Hence g(n, m) is periodic in the variable m with period 1.

The frequency-domain relationships can be obtained by combining the results of the interpolation and decimation processes. Thus the spectrum at the output of the linear filter with impulse response h(l) is

$$V(\omega_{\nu}) = H(\omega_{\nu})X(\omega_{\nu}I)$$

=
$$\begin{cases} IX(\omega_{\nu}I), & 0 \le |\omega_{\nu}| \le \min(\pi/D, \pi/I) \\ 0, & \text{otherwise} \end{cases}$$
(10.4.10)

The spectrum of the output sequence y(m), obtained by decimating the sequence v(n) by a factor of D, is

$$Y(\omega_{y}) = \frac{1}{D} \sum_{k=0}^{D-1} V\left(\frac{\omega_{y} - 2\pi k}{D}\right)$$
(10.4.11)

where $\omega_y = D\omega_v$. Since the linear filter prevents aliasing as implied by (10.4.10), the spectrum of the output sequence given by (10.4.11) reduces to

$$Y(\omega_{y}) = \begin{cases} \frac{l}{D} X\left(\frac{\omega_{y}}{D}\right), & 0 \le |\omega_{y}| \le \min\left(\pi, \frac{\pi D}{l}\right) \\ 0, & \text{otherwise} \end{cases}$$
(10.4.12)

UNIT V DIGITAL SIGNAL PROCESSOR

INTRODUCTION

A digital signal processor (DSP) is a specialized microprocessor (or a SIP block), with architecture optimized for the operational needs of digital signal processing.

The goal of DSP is usually to measure, filter or compress continuous real-world analog signals. Most general-purpose microprocessors can also execute digital signal processing algorithms successfully, but may not be able to keep up with such processing continuously in real-time. Also, dedicated DSPs usually have better power efficiency, thus they are more suitable in portable devices such as mobile phones because of power consumption constraints. DSPs often use special memory architectures that are able to fetch multiple data or instructions at the same time.

Circular Buffering

Digital Signal Processors are designed to quickly carry out FIR filters and similar techniques. To understand the *hardware*, we must first understand the *algorithms*. In this section we will make a detailed list of the steps needed to implement an FIR filter. In the next section we will see how DSPs are designed to perform these steps as efficiently as possible.

To start, we need to distinguish between **off-line processing** and **real-time processing**. In offline processing, the *entire* input signal resides in the computer at the same time. For example, a geophysicist might use a seismometer to record the ground movement during an earthquake. After the shaking is over, the information may be read into a computer and analyzed in some way. Another example of off-line processing is medical imaging, such as computed tomography and MRI. The data set is acquired while the patient is inside the machine, but the image reconstruction may be delayed until a later time. The key point is that *all* of the information is simultaneously available to the processing program. This is common in scientific research and engineering, but not in consumer products. Off-line processing is the realm of personal computers and mainframes.

In real-time processing, the output signal is produced at the same time that the input signal is being acquired. For example, this is needed in telephone communication, hearing aids, and radar . These applications must have the information immediately available, although it can be delayed by a short amount. For instance, a 10 millisecond delay in a telephone call cannot be detected by the speaker or listener. Likewise, it makes no difference if a radar signal is delayed by a few seconds before being displayed to the operator. Real-time applications input a sample, perform the algorithm, and output a sample, over-and-over. Alternatively, they may input a group

of samples, perform the algorithm, and output a group of samples. This is the world of Digital Signal Processors.



Now look at Fig. 28-2 and imagine that this is an FIR filter being implemented in *real-time*. To calculate the output sample, we must have access to a certain number of the most recent samples from the input. For example, suppose we use eight coefficients in this filter, a_0, a_1, \ldots, a_7 . This means we must know the value of the eight most recent samples from the input signal, x[n], x[n-1], $\ldots x[n-7]$. These eight samples must be stored in memory and continually updated as new samples are acquired. What is the best way to manage these stored samples? The answer is **circular buffering**.



Figure 28-3 illustrates an eight sample circular buffer. We have placed this circular buffer in eight consecutive memory locations, 20041 to 20048. Figure (a) shows how the eight samples from the input might be stored at one particular instant in time, while (b) shows the changes after the next sample is acquired. The idea of circular buffering is that the end of this linear array is connected to its beginning; memory location 20041 is viewed as being next to 20048, just as 20044 is next to 20045. You keep track of the array by a **pointer** (a variable whose value is an *address*) that indicates where the most recent sample resides.

For instance, in (a) the pointer contains the address 20044, while in (b) it contains 20045. When a new sample is acquired, it replaces the oldest sample in the array, and the pointer is moved one address ahead. Circular buffers are efficient because only one value needs to be changed when a new sample is acquired.

Four parameters are needed to manage a circular buffer. First, there must be a pointer that indicates the start of the circular buffer in memory (in this example, 20041). Second, there must be a pointer indicating the end of the array (e.g., 20048), or a variable that holds its length (e.g., 8). Third, the step size of the memory addressing must be specified. In Fig. 28-3 the step size is *one*, for example: address 20043 contains one sample, address 20044 contains the next sample, and so on. This is frequently not the case. For instance, the addressing may refer to bytes, and each sample may require two or four bytes to hold its value. In these cases, the step size would need to be two or four, respectively.

These three values define the size and configuration of the circular buffer, and will not change during the program operation. The fourth value, the pointer to the most recent sample, must be

modified as each new sample is acquired. In other words, there must be program logic that controls how this fourth value is updated based on the value of the first three values. While this logic is quite simple, it must be very fast. This is the whole point of this discussion; DSPs should be optimized at managing circular buffers to achieve the highest possible execution speed.

As an aside, circular buffering is also useful in *off-line* processing. Consider a program where both the input and the output signals are completely contained in memory. Circular buffering isn't needed for a convolution calculation, because every sample can be immediately accessed. However, many algorithms are implemented in *stages*, with an intermediate signal being created between each stage. For instance, a recursive filter carried out as a series of biquads operates in this way. The brute force method is to store the entire length of each intermediate signal in memory. Circular buffering provides another option: store only those intermediate samples needed for the calculation at hand. This reduces the required amount of memory, at the expense of a more complicated algorithm. The important idea is that circular buffers are *useful* for off-line processing, but *critical* for real-time applications.

Now we can look at the steps needed to implement an FIR filter using circular buffers for both the input signal and the coefficients. This list may seem trivial and overexamined- it's not! The efficient handling of these individual tasks is what separates a DSP from a traditional microprocessor. For each new sample, all the following steps need to be taken:

TABLE 28-1 FIR filter steps.	 Obtain a sample with the ADC; generate an interrupt Detect and manage the interrupt Move the sample into the input signal's circular buffer Update the pointer for the input signal's circular buffer Zero the accumulator Control the loop through each of the coefficients Fetch the coefficient from the coefficient's circular buffer Update the pointer for the input signal's circular buffer Fetch the sample from the coefficient's circular buffer Fetch the sample from the input signal's circular buffer Update the pointer for the input signal's circular buffer Update the pointer for the input signal's circular buffer Update the pointer for the input signal's circular buffer Update the pointer for the input signal's circular buffer Update the pointer for the input signal's circular buffer Multiply the coefficient by the sample Add the product to the accumulator Move the output sample (accumulator) to a holding buffer Move the output sample from the holding buffer to the DAC
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The goal is to make these steps execute quickly. Since steps 6-12 will be repeated many times (once for each coefficient in the filter), special attention must be given to these operations. Traditional microprocessors must generally carry out these 14 steps in *serial* (one after another), while DSPs are designed to perform them in *parallel*. In some cases, all of the operations within the loop (steps 6-12) can be completed in a *single clock cycle*. Let's look at the internal architecture that allows this magnificent performance.

ARCHITECTURE OF THE DIGITAL SIGNAL PROCESSOR

One of the biggest bottlenecks in executing DSP algorithms is transferring information to and from memory. This includes *data*, such as samples from the input signal and the filter coefficients, as well as *program instructions*, the binary codes that go into the program sequencer. For example, suppose we need to multiply two numbers that reside somewhere in memory. To do this, we must fetch three binary values from memory, the numbers to be multiplied, plus the program instruction describing what to do.

Figure 28-4a shows how this seemingly simple task is done in a traditional microprocessor. This is often called a **Von Neumann architecture**, after the brilliant American mathematician John Von Neumann (1903-1957). Von Neumann guided the mathematics of many important discoveries of the early twentieth century. His many achievements include: developing the concept of a stored program computer, formalizing the mathematics of quantum mechanics, and work on the atomic bomb. If it was new and exciting, Von Neumann was there!

As shown in (a), a Von Neumann architecture contains a single memory and a single bus for transferring data into and out of the central processing unit (CPU). Multiplying two numbers requires at least three clock cycles, one to transfer each of the three numbers over the bus from the memory to the CPU. We don't count the time to transfer the result back to memory, because we assume that it remains in the CPU for additional manipulation (such as the sum of products in an FIR filter). The Von Neumann design is quite satisfactory when you are content to execute all of the required tasks in serial. In fact, most computers today are of the Von Neumann design. We only need other architectures when very fast processing is required, and we are willing to pay the price of increased complexity.

This leads us to the **Harvard architecture**, shown in (b). This is named for the work done at Harvard University in the 1940s under the leadership of Howard Aiken (1900-1973). As shown in this illustration, Aiken insisted on separate memories for data and program instructions, with separate buses for each. Since the buses operate independently, program instructions and data can be fetched at the same time, improving the speed over the single bus design. Most present day DSPs use this dual bus architecture.

Figure (c) illustrates the next level of sophistication, the **Super Harvard Architecture**. This term was coined by Analog Devices to describe the internal operation of their ADSP-2106x and new ADSP -211 xx families of Digital Signal Processors . These are called **SHARC** ® DSPs, a contraction of the longer term, <u>Super Harvard ARC</u>hitecture . The idea is to build upon the Harvard architecture by adding features to improve the throughput. While the SHARC DSPs are optimized in dozens of ways, two areas are important enough to be included in Fig. 28-4c: an *instruction cache*, and an *I/O controller*.

First, let's look at how the instruction cache improves the performance of the Harvard architecture. A handicap of the basic Harvard design is that the data memory bus is busier than the program memory bus. When two numbers are multiplied, two binary values (the numbers) must be passed over the data memory bus, while only one binary value (the program instruction) is passed over the program memory bus. To improve upon this situation, we start by relocating part of the "data" to program memory. For instance, we might place the filter coefficients in program memory, while keeping the input signal in data memory. (This relocated data is called " secondary data" in the illustration). At first glance, this doesn't seem to help the situation; now we must transfer one value over the data memory bus (the input signal sample), but two values over the program memory bus (the program instruction and the coefficient). In fact, if we were executing random instructions, this situation would be no better at all.

However, DSP algorithms generally spend most of their execution time in loops, such as instructions 6-12 of Table 28-1. This means that the same set of program instructions will continually pass from program memory to the CPU. The Super Harvard architecture takes advantage of this situation by including an **instruction cache** in the CPU. This is a small memory that contains about 32 of the most recent program instructions. The first time through a loop, the program instructions must be passed over the program memory bus. This results in slower operation because of the conflict with the coefficients that must also be fetched along this path. However, on additional executions of the loop, the program instructions can be pulled from the instruction cache. This means that all of the memory to CPU information transfers can be accomplished in a single cycle: the sample from the input signal comes over the data memory bus, the coefficient comes over the program memory bus, and the program instruction comes from the instruction cache. In the jargon of the field, this efficient transfer of data is called a *high memory-access bandwidth*.

Figure 28-5 presents a more detailed view of the SHARC architecture, showing the **I/O controller** connected to data memory. This is how the signals enter and exit the system. For instance, the SHARC DSPs provides both serial and parallel communications ports. These are extremely high speed connections. For example, at a 40 MHz clock speed, there are two serial ports that operate at 40 Mbits/second each, while six parallel ports each provide a 40 Mbytes/second data transfer. When all six parallel ports are used together, the data transfer rate is an incredible 240 Mbytes/second.



This is fast enough to transfer the entire text of this book in only 2 milliseconds! Just as important, dedicated hardware allows these data streams to be transferred directly into memory (Direct Memory Access, or DMA), without having to pass through the CPU's registers. In other words, tasks 1 & 14 on our list happen independently and simultaneously with the other tasks; no cycles are stolen from the CPU. The main buses (program memory bus and data memory bus) are also accessible from outside the chip, providing an additional interface to off-chip memory and peripherals. This allows the SHARC DSPs to use a four Gigaword (16 Gbyte) memory, accessible at 40 Mwords/second (160 Mbytes/second), for 32 bit data. Wow!

This type of high speed I/O is a key characteristic of DSPs. The overriding goal is to move the data in, perform the math, and move the data out before the next sample is available. Everything else is secondary. Some DSPs have on-board analog-to-digital and digital-to-analog converters, a feature called **mixed signal**. However, all DSPs can interface with external converters through serial or parallel ports.

Now let's look inside the CPU. At the top of the diagram are two blocks labeled **Data Address Generator** (DAG), one for each of the two memories. These control the addresses sent to the program and data memories, specifying where the information is to be read from or written to. In simpler microprocessors this task is handled as an inherent part of the program sequencer, and is quite transparent to the programmer. However, DSPs are designed to operate with *circular buffers*, and benefit from the extra hardware to manage them efficiently. This avoids needing to use precious CPU clock cycles to keep track of how the data are stored. For instance, in the

SHARC DSPs, each of the two DAGs can control *eight* circular buffers. This means that each DAG holds 32 variables (4 per buffer), plus the required logic.

Why so many circular buffers? Some DSP algorithms are best carried out in stages. For instance, IIR filters are more stable if implemented as a cascade of biquads (a stage containing two poles and up to two zeros). Multiple stages require multiple circular buffers for the fastest operation. The DAGs in the SHARC DSPs are also designed to efficiently carry out the *Fast Fourier transform*. In this mode, the DAGs are configured to generate **bit-reversed addresses** into the circular buffers, a necessary part of the FFT algorithm. In addition, an abundance of circular buffers greatly simplifies DSP code generation- both for the human programmer as well as high-level language compilers, such as C.

The data register section of the CPU is used in the same way as in traditional microprocessors. In the ADSP-2106x SHARC DSPs, there are 16 general purpose registers of 40 bits each. These can hold intermediate calculations, prepare data for the math processor, serve as a buffer for data transfer, hold flags for program control, and so on. If needed, these registers can also be used to control loops and counters; however, the SHARC DSPs have extra hardware registers to carry out many of these functions.

The math processing is broken into three sections, a **multiplier**, an **arithmetic logic unit** (**ALU**), and a **barrel shifter**. The multiplier takes the values from two registers, multiplies them, and places the result into another register. The ALU performs addition, subtraction, absolute value, logical operations (AND, OR, XOR, NOT), conversion between fixed and floating point formats, and similar functions. Elementary binary operations are carried out by the barrel shifter, such as shifting, rotating, extracting and depositing segments, and so on. A powerful feature of the SHARC family is that the multiplier and the ALU can be accessed in parallel. In a single clock cycle, data from registers 0-7 can be passed to the multiplier, data from registers 8-15 can be passed to the ALU, and the two results returned to any of the 16 registers.

There are also many important features of the SHARC family architecture that aren't shown in this simplified illustration. For instance, an 80 bit accumulator is built into the multiplier to reduce the round-off error associated with multiple fixed-point math operations. Another interesting



feature is the use of **shadow registers** for all the CPU's key registers. These are duplicate registers that can be switched with their counterparts in a single clock cycle. They are used for *fast context switching*, the ability to handle interrupts quickly. When an interrupt occurs in traditional microprocessors, all the internal data must be saved before the interrupt can be handled. This usually involves pushing all of the occupied registers onto the stack, one at a time. In comparison, an interrupt in the SHARC family is handled by moving the internal data into the shadow registers in a *single clock cycle*. When the interrupt routine is completed, the registers are just as quickly restored. This feature allows step 4 on our list (managing the sample-ready interrupt) to be handled very quickly and efficiently.

Now we come to the critical performance of the architecture, how many of the operations within the loop (steps 6-12 of Table 28-1) can be carried out at the same time. Because of its highly parallel nature, the SHARC DSP can simultaneously carry out *all* of these tasks. Specifically, within a single clock cycle, it can perform a multiply (step 11), an addition (step 12), two data moves (steps 7 and 9), update two circular buffer pointers (steps 8 and 10), and control the loop (step 6). There will be extra clock cycles associated with beginning and ending the loop (steps 3, 4, 5 and 13, plus moving initial values into place); however, these tasks are also handled very efficiently. If the loop is executed more than a few times, this overhead will be negligible. As an example, suppose you write an efficient FIR filter program using 100 coefficients. You can

expect it to require about 105 to 110 clock cycles per sample to execute (i.e., 100 coefficient loops plus overhead). This is very impressive; a traditional microprocessor requires many thousands of clock cycles for this algorithm.

FIXED VERSUS FLOATING POINT

Digital Signal Processing can be divided into two categories, **fixed point** and **floating point**. These refer to the format used to store and manipulate numbers within the devices.

Fixed point DSPs usually represent each number with a minimum of 16 bits, although a different length can be used. For instance, Motorola manufactures a family of fixed point DSPs that use 24 bits. There are four common ways that these $2^{16} = 65536$ possible bit patterns can represent a number.

In **unsigned integer**, the stored number can take on any integer value from 0 to 65,535. Similarly, **signed integer** uses two's complement to make the range include negative numbers, from -32,768 to 32,767. With **unsigned fraction** notation, the 65,536 levels are spread uniformly between 0 and 1. Lastly, the **signed fraction** format allows negative numbers, equally spaced between -1 and 1.

In comparison, floating point DSPs typically use a minimum of 32 bits to store each value. This results in many more bit patterns than for fixed point, $2^{32} = 4,294,967,296$ to be exact. A key feature of floating point notation is that the represented numbers are *not* uniformly spaced. In the most common format (ANSI/IEEE Std. 754-1985), the largest and smallest numbers are $\pm 3.4 \times 10^{38}$ and 1.2 10^{-38} , respectively.

The represented values are unequally spaced between these two extremes, such that the gap between any two numbers is about ten-million times smaller than the value of the numbers. This is important because it places large gaps between large numbers, but small gaps between small numbers. Floating point notation is discussed in more detail in Chapter 4.

All floating point DSPs can also handle fixed point numbers, a necessity to implement counters, loops, and signals coming from the ADC and going to the DAC. However, this doesn't mean that fixed point math will be carried out as quickly as the floating point operations; it depends on the internal architecture. For instance, the SHARC DSPs are optimized for both floating point and fixed point operations, and executes them with equal efficiency. For this reason, the SHARC devices are often referred to as "32-bit DSPs," rather than just "Floating Point."

Figure 28-6 illustrates the primary trade-offs between fixed and floating point DSPs. In Chapter 3 we stressed that fixed point arithmetic is much